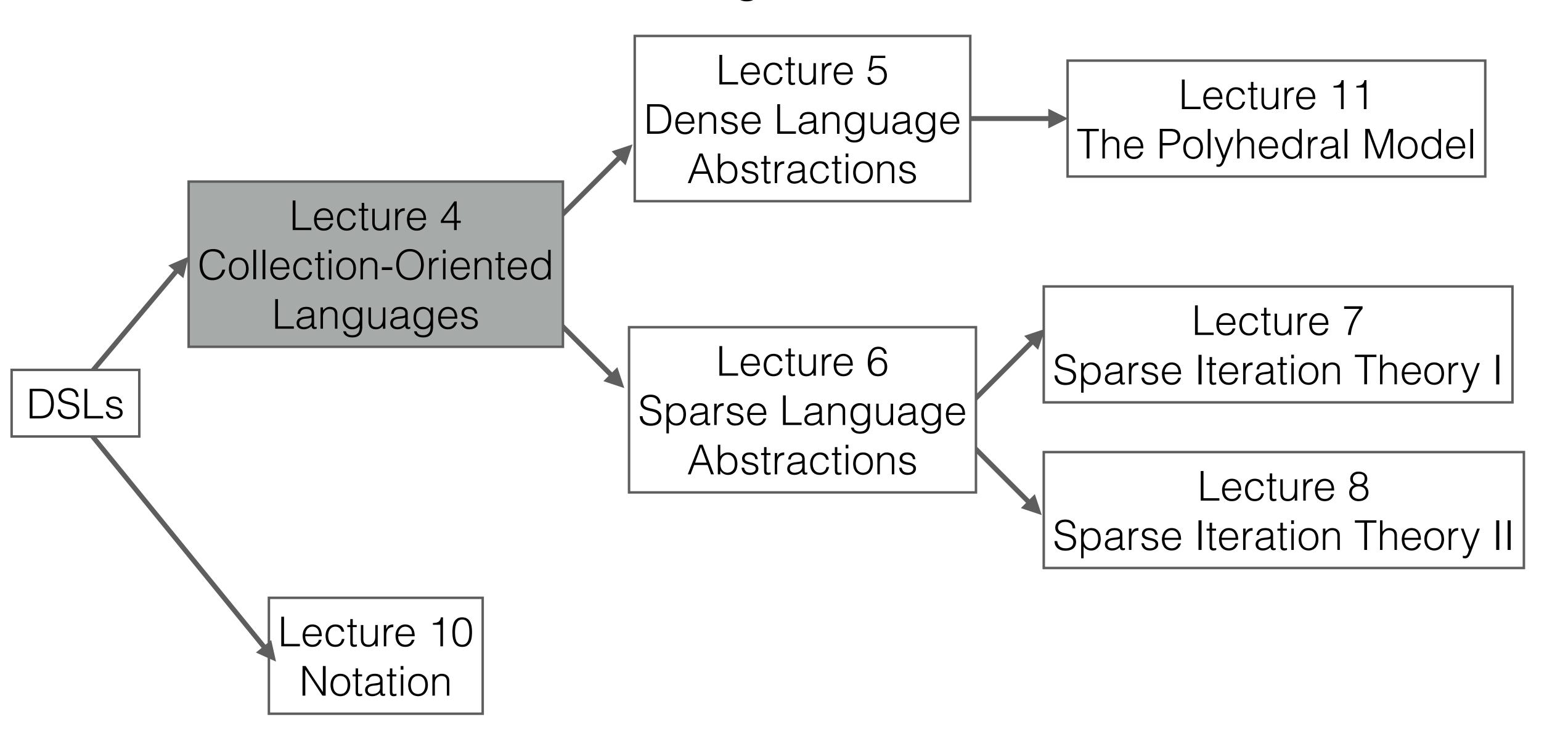
Lecture 4 - Collection-Oriented Languages

Stanford CS343D (Fall 2020) Fred Kjolstad and Pat Hanrahan

Overview of lectures in the coming weeks



Languages are tools for thought

"By relieving the brain of all unnecessary work, a good notation sets it free to concentrate on the more advanced problems, and in effect increases the mental power of the race."

— Alfred N. Whitehead

Collection-Oriented languages are an important subclass of DSLs as discussed in this course

DSLs

Collection-Oriented Languages

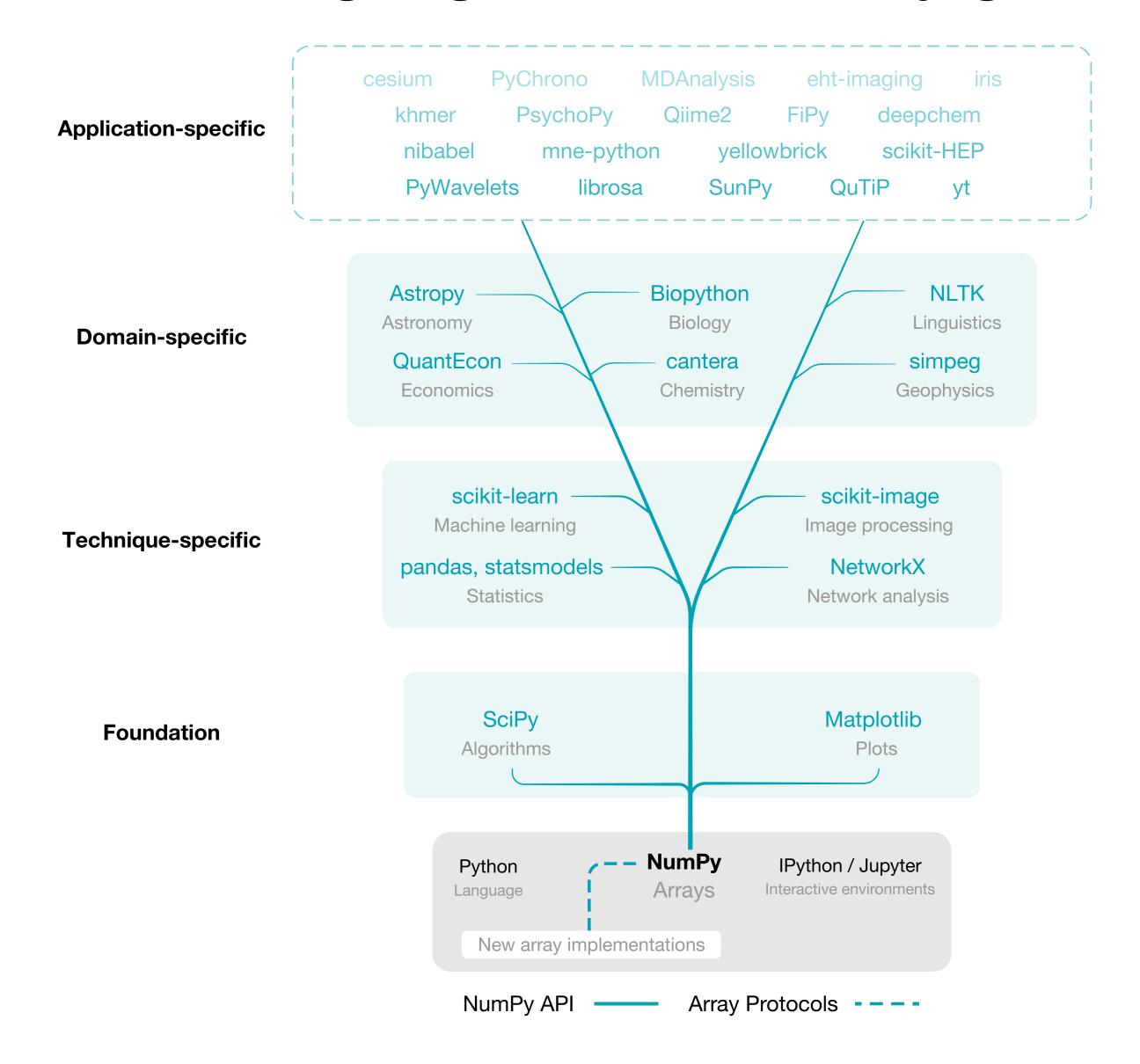
Economy of scale in notation and execution

$$C = A \bowtie B$$

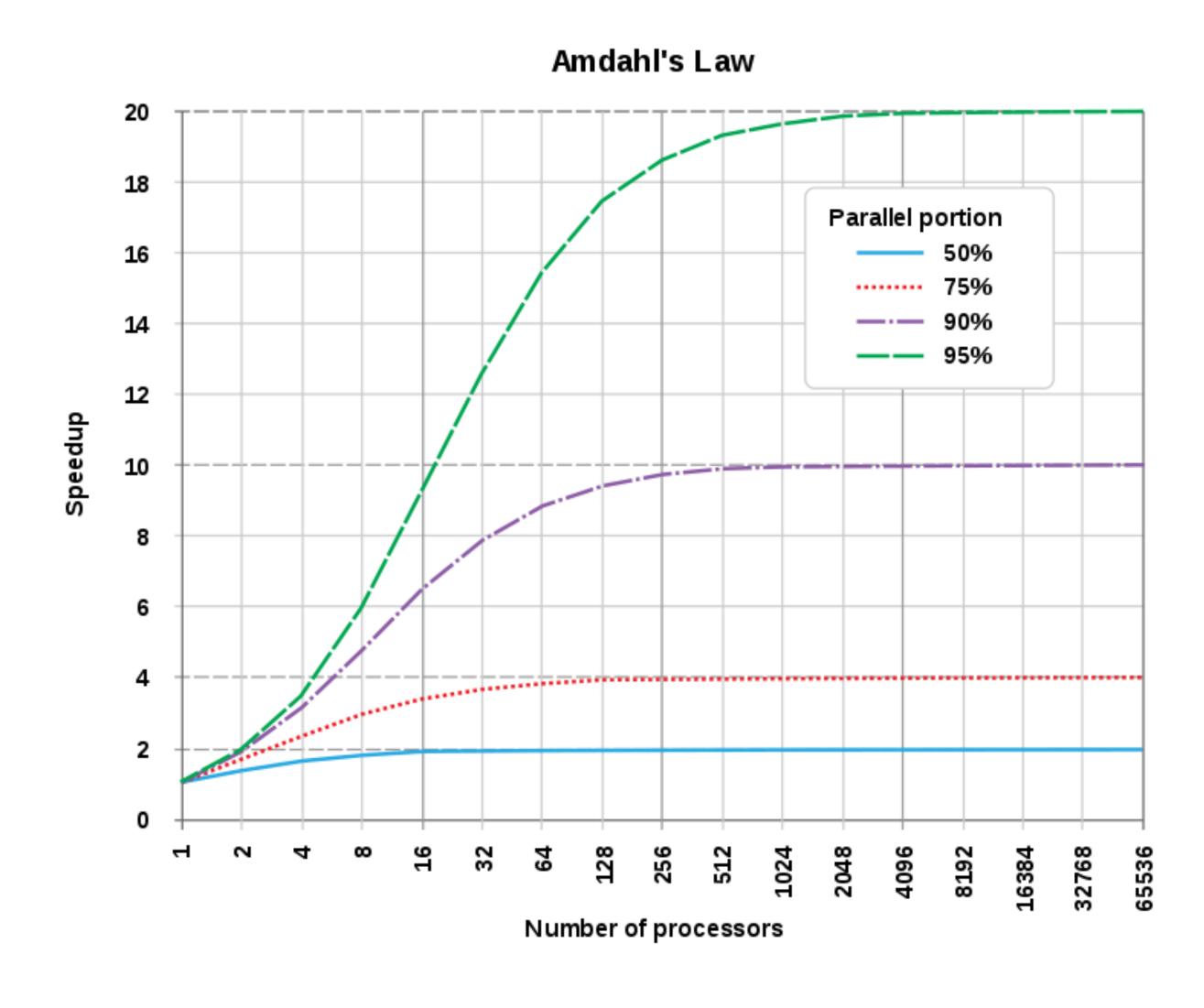
$$c = Ab$$

How many operations?

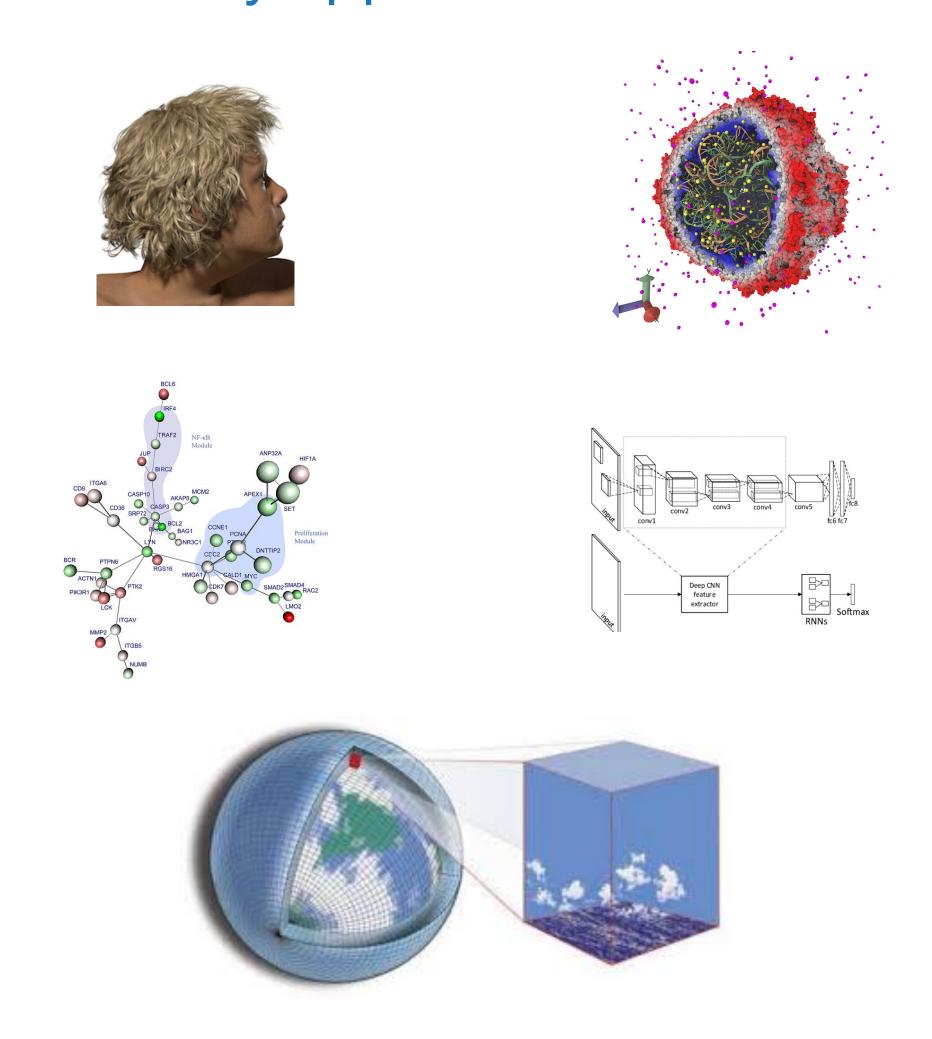
Collection-oriented languages are relatively general



We need collections for performance due to Amdahl's law



But many applications are data-rich



Plot from Wikipedia

Can programming be liberated from the von Neumann model?

Imperative Form

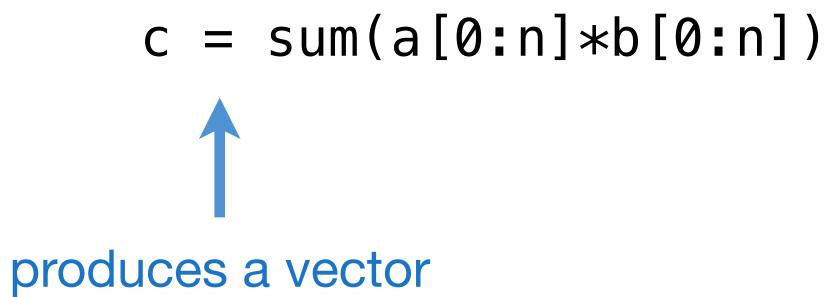
c := 0
for i := 1 step 1 until n do
 c := c + a[i] x b[i]

poor world of statements

rich world of expressions

transfers one scalar value to memory:
von Neumann bottleneck in software
the assignment transfers one value to memory

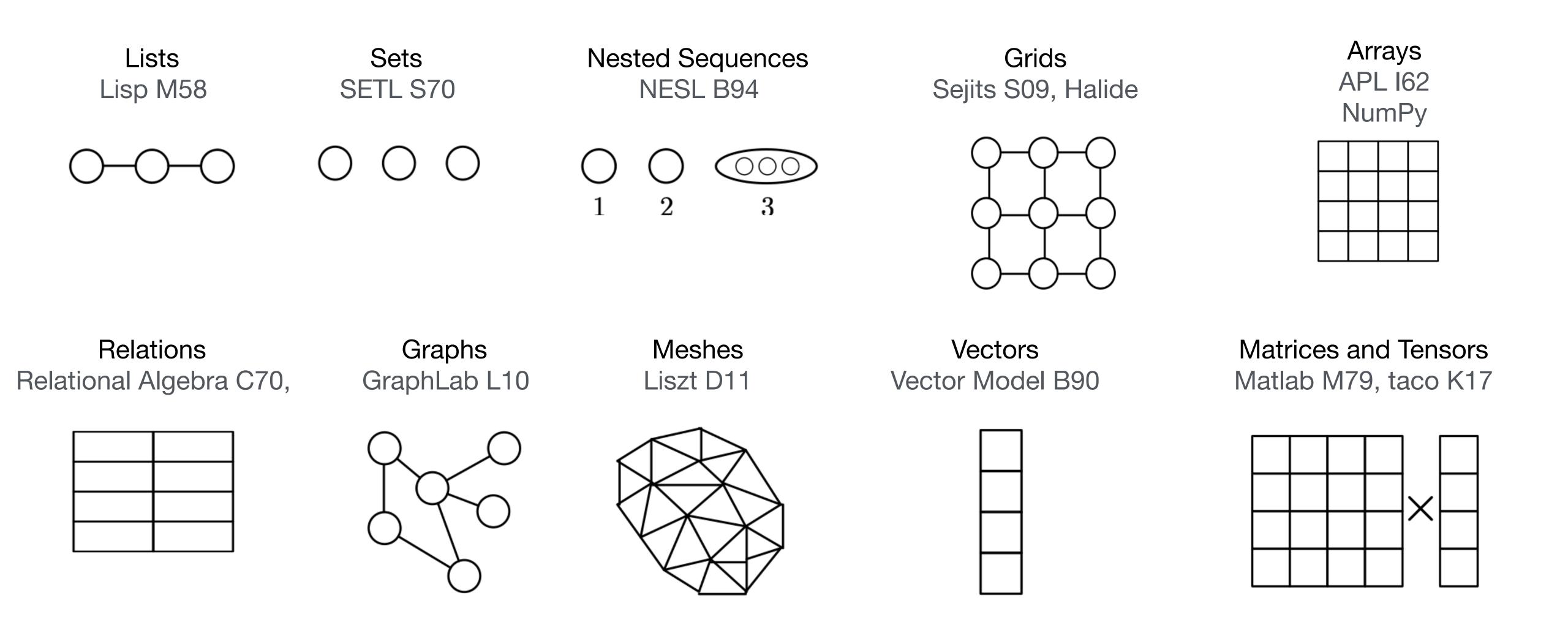
Functional Form



Collection-oriented operations let us operate on collections as a whole

- A record-at-a-time user interface forces the programmer to do manual query optimization, and this is often hard.
- Set-a-time languages are good, regardless of the data model, since they offer much improved physical data independence.
- The programming language community has long recognized that aggregate data structures and general operations on them give great flexibility to programmers and language implementors.

Collection-Oriented Languages



A collection-oriented programming model provides collective operations on some collection/abstract data structure

Features of collections

- Ordering: unordered, sequence, or grid-ordered?
- Regularity: Can the collection represent irregularity/sparsity?
- Nesting: nested or flat collections?
- Random-access: can individual elements be accessed?

The APL Programming Language

n	←	4	5	6	7			
	+4 9	10) [11				

i.e., mkArray

4 is broadcast across each n

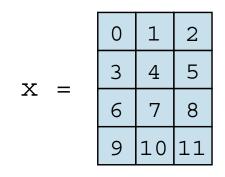
element-wise multiplication (14 makes the array [1,2,3,4])

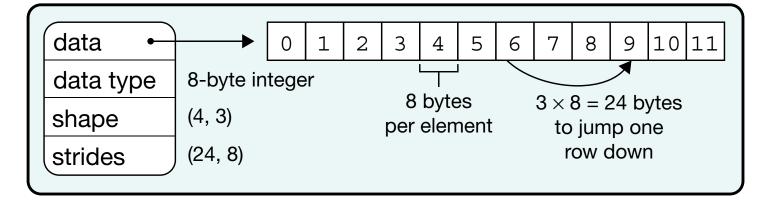
$$\sum_{i=0}^{n} n_i$$

$$\sum_{i=1}^{4} (i+3)$$

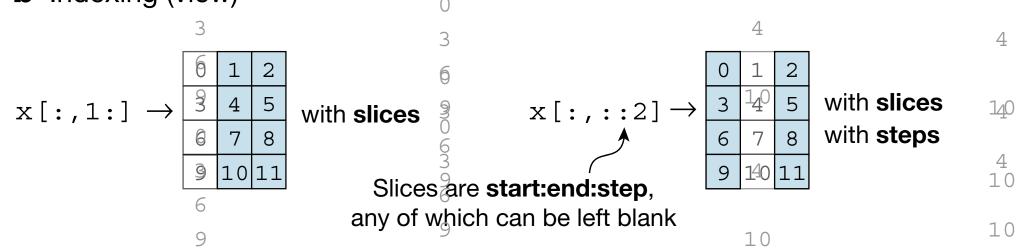
Array Programming with NumPy

a Data structure

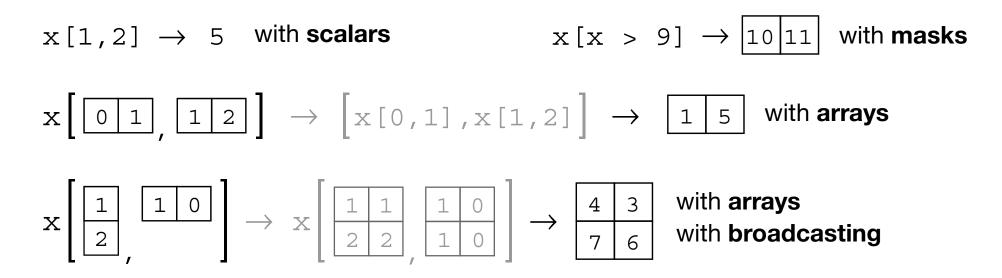




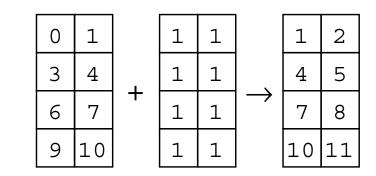
b Indexing (view)



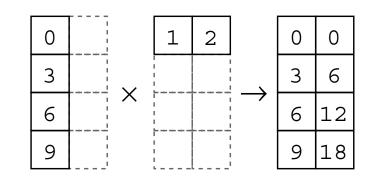
c Indexing (copy)



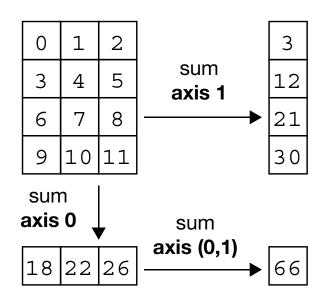
d Vectorization



e Broadcasting



f Reduction



The SETL Language

Sets Tuples Functions

(O,O,O)

SETL Set Former Notation

Notation

$$\{x \in s \mid C(x)\}$$

$$\{x \in \{1,5,10,32\} \mid x \text{ lt } 10\} \rightarrow \{1,5\}$$

$$\{e(x), x \in s \mid C(x)\}$$

$$\{i*i, i \in \{1,3,5\}\} \rightarrow \{1,9,25\}$$

$$\{e(x), \min \leq i \leq \max \mid C(i)\}$$

$$\{i*2-1,1 \le i \le 5\} \rightarrow \{1,3,5,7,9\}$$

$$[op: x \in s \mid C(x)]e(x)$$

$$[+:x \in \{1,2,3\}](x*x) \to 14$$

$$\forall x \in s \mid C(x)$$

$$\forall x \in 1,2,4 \mid (x//2) \text{ eq } 1 \to \mathbf{f}$$

$$[+:x \in s_1, y \in s_2]\{ < x, y > \}$$

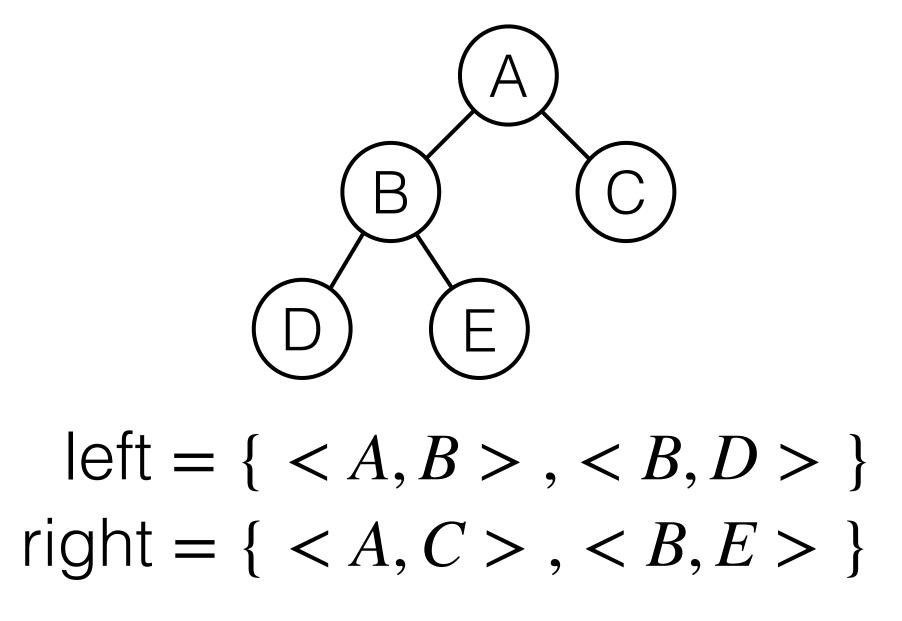
$$[+:x \in \{1,2\}, y \in \{a,b\}]\{\langle x,y \rangle\} \rightarrow \{\langle 1,a \rangle,\langle 1,b \rangle,\langle 2,a \rangle,\langle 2,b \rangle\}$$

Standard set operations like union, intersection, and set difference are also supported

SETL Table Functions

$$f = \{ < 1,1 > , < 2,4 > , < 3,9 > \}$$

 $f(2) \to 4$
 $f + \{ < 2,5 > \} \to \{ < 1,1 > , < 2,5 > , < 3,9 > \}$



Relational Algebra

employees

name	id	department
Harry	3245	CS
Sally	7264	EE
George	1379	CS
Mary	1733	ME
Rita	2357	CS

departments

department	manager
CS	George
EE	Mary

Projection (Π)

 $\Pi_{name,department}$ employees

Name	Department
Harry	CS
Sally	EE
George	CS
Mary	ME
Rita	CS

Selection (σ)

 $\sigma_{department=CS}$ employees

Name	ID	Department
Harry	3245	CS
George	1379	CS
Rita	2357	CS

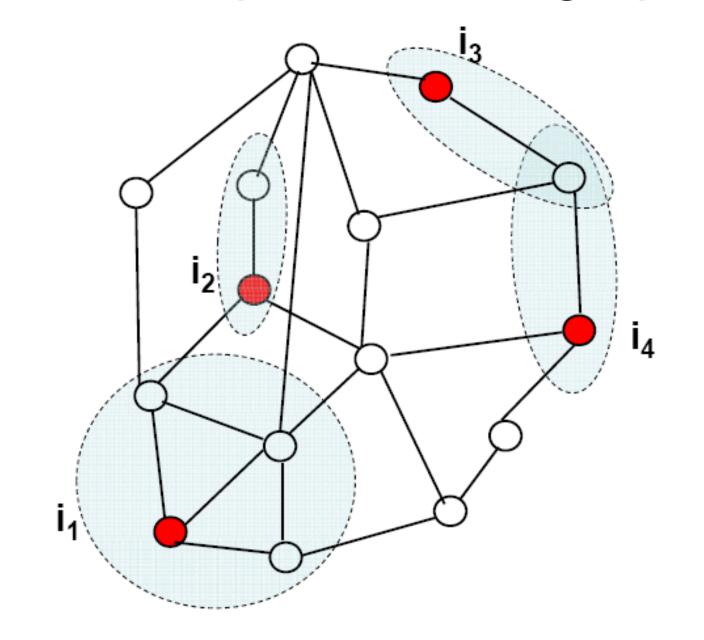
Natural join (⋈)

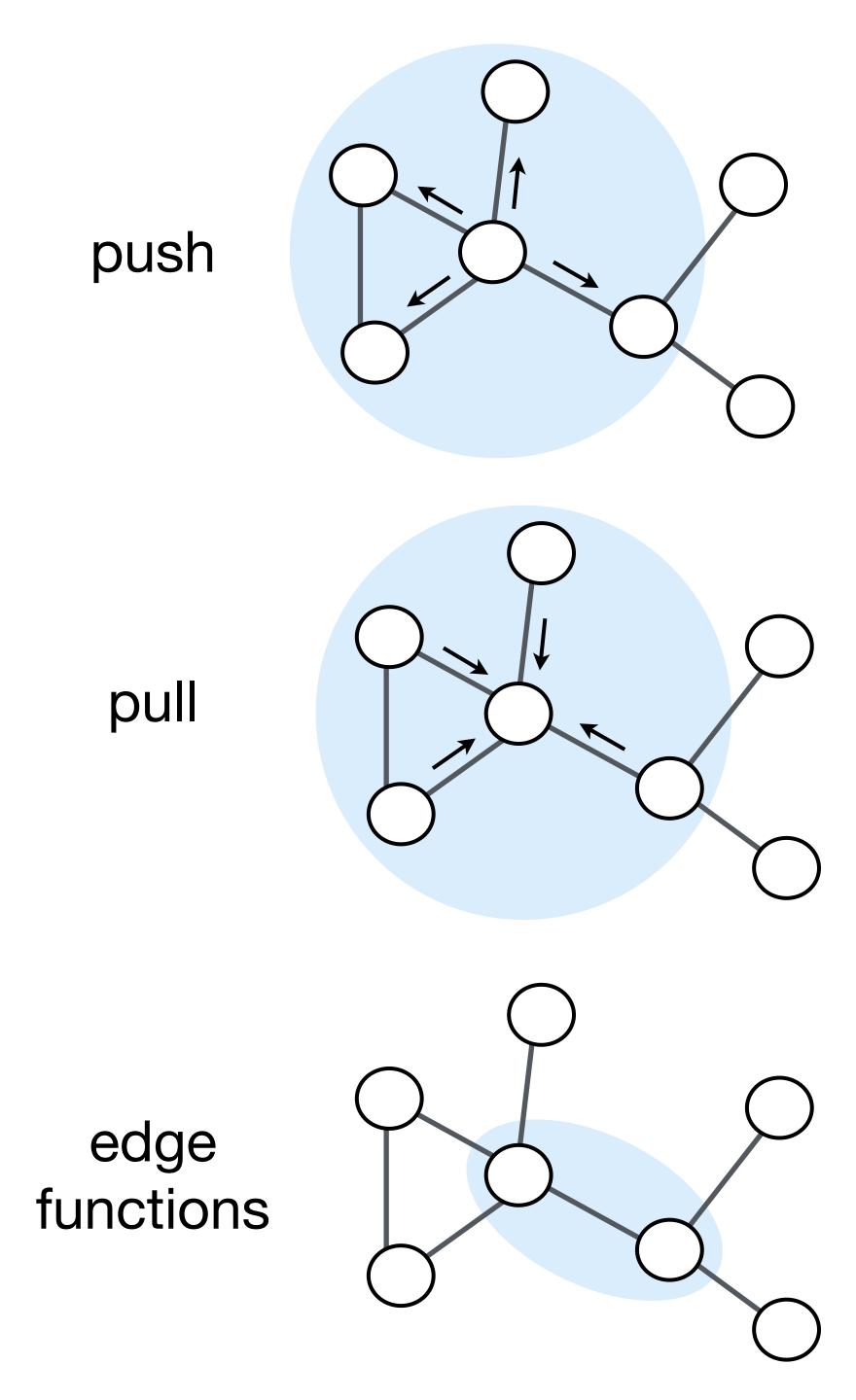
employees M departments

name	id	department	manager
Harry	3245	CS	George
Sally	7264	EE	Mary
George	1379	CS	George
Rita	2357	CS	George

Graph operations

Simultaneous operations on different parts of the graph





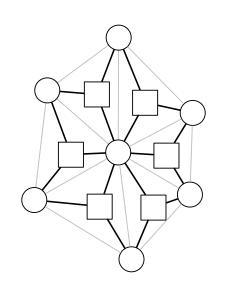
Relations, Graphs, and Algebra: No glove fits all

Relations

Names	City	Age
Peter	Boston	54
Mary	San Fransisco	35
Paul	New York	23
Adam	Seattle	84
Hilde	Boston	19
Bob	Chicago	76
Sam	Portland	32
Angela	Los Angeles	62

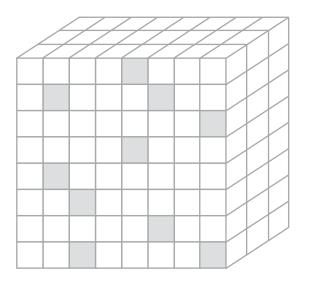
Ideal for combining data to form systems

Graphs



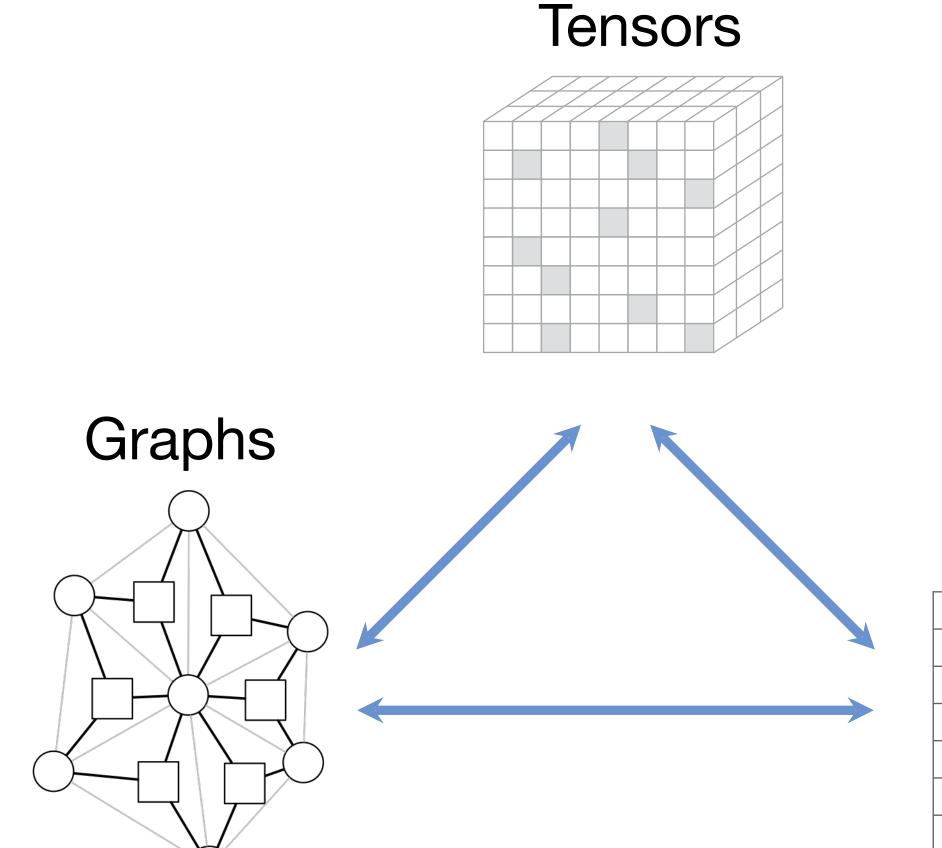
Ideal for local operations

Tensors



Ideal for global operations

It is critical to be able to compose languages and abstractions

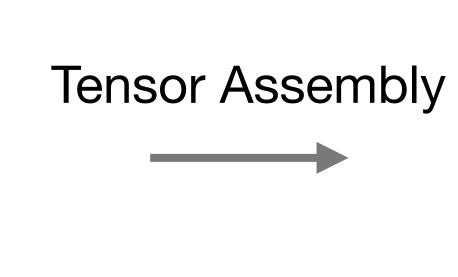


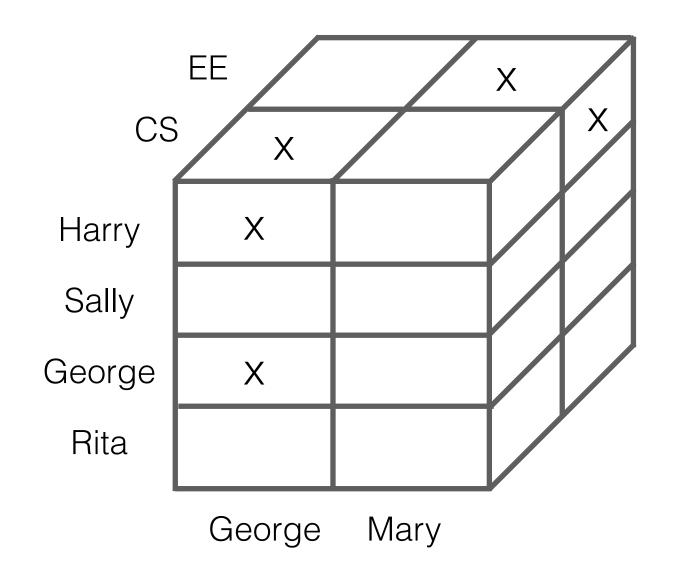
Relations

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Hilde	Boston	19
Bob	Chicago	76
Sam	Portland	32
Angela	Los Angeles	62

Example: Relations and Tensors

department	manager
CS	George
EE	Mary
CS	George
CS	George
	CS EE CS



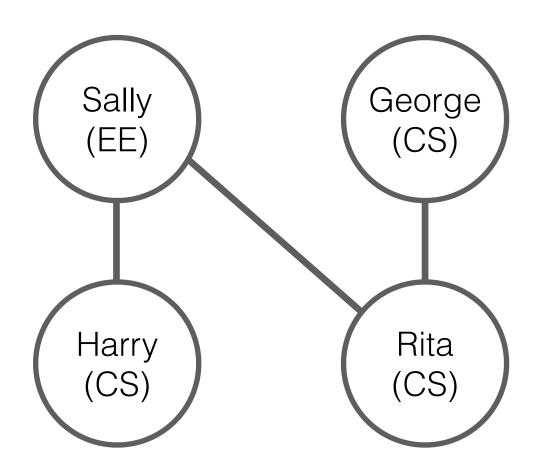


Example: Relations and Graphs

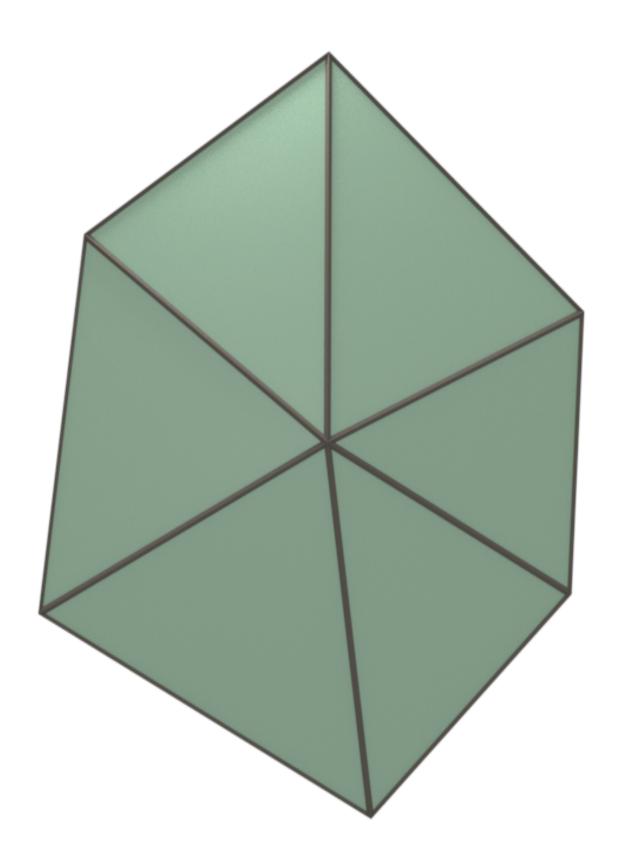
name	department
Harry	CS
Sally	EE
George	CS
Rita	CS

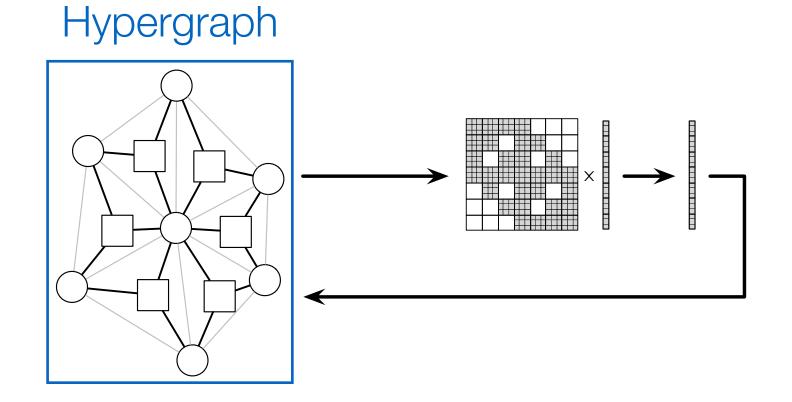


name1	name2
Harry	Sally
Sally	Harry
George	Rita
Rita	George
Sally	Rita
Rita	Sally

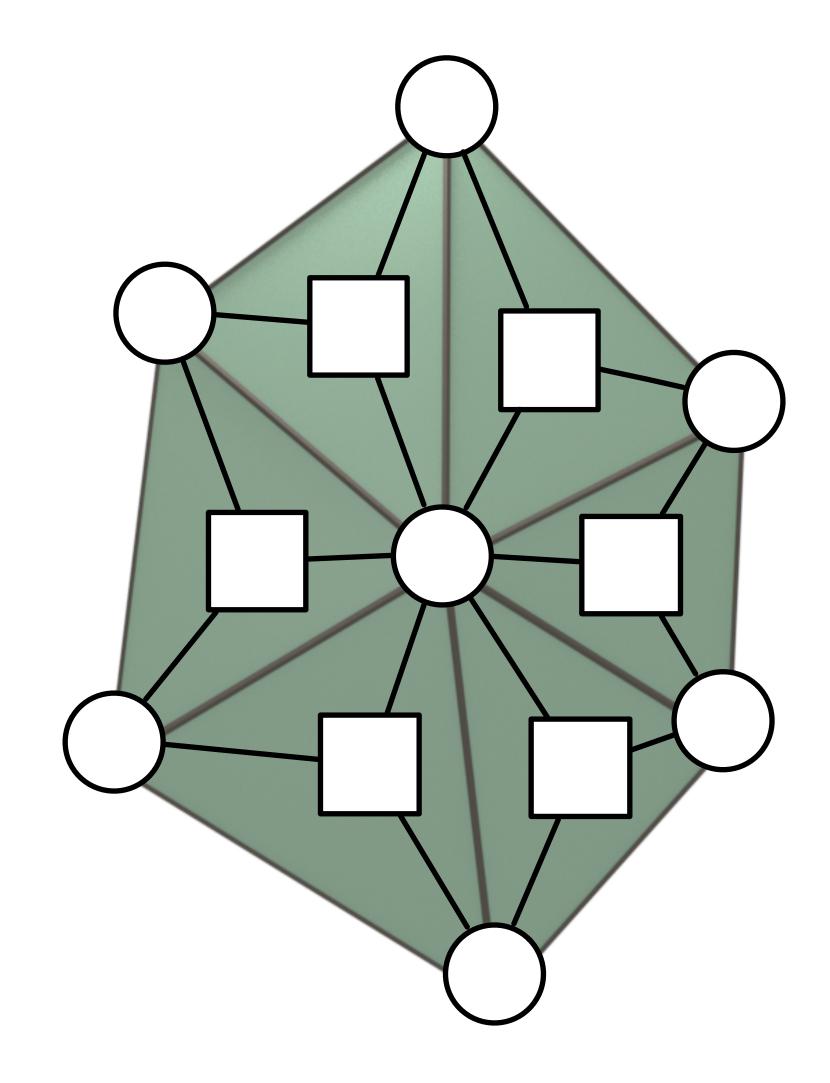


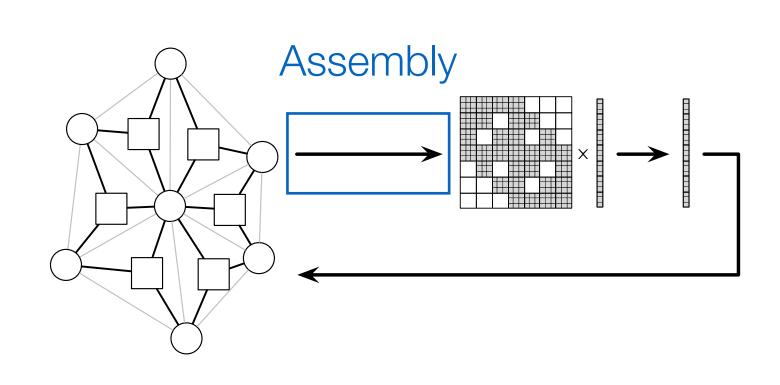
Example: Graphs and Tensors (Simit)





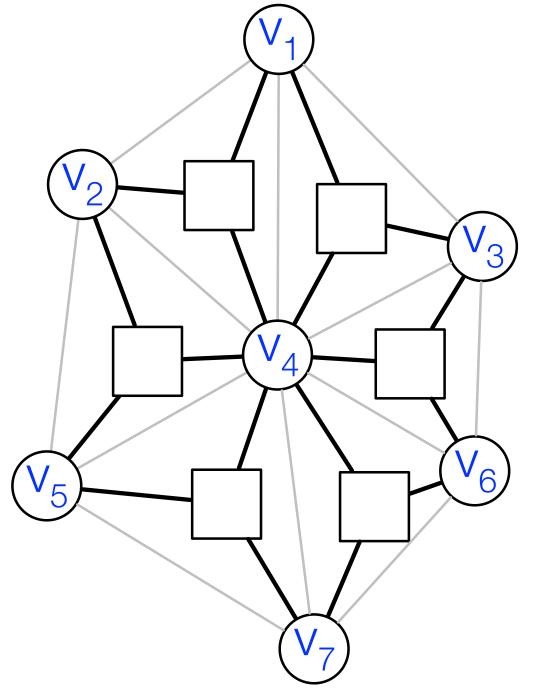
```
element Vertex
 x : vector[3](float); % position
 v : vector[3](float); % velocity
 fe : vector[3](float); % external force
end
element Triangle
                     % shear modulus
 u : float;
 l : float;
                     % lame's first parameter
 W : float;
                        % volume
    : matrix[3,3](float); % strain-displacement
end
% graph vertices and triangle hyperedges
               : set{Vertex};
extern verts
extern triangles : set{Triangle}(verts, verts, verts);
```

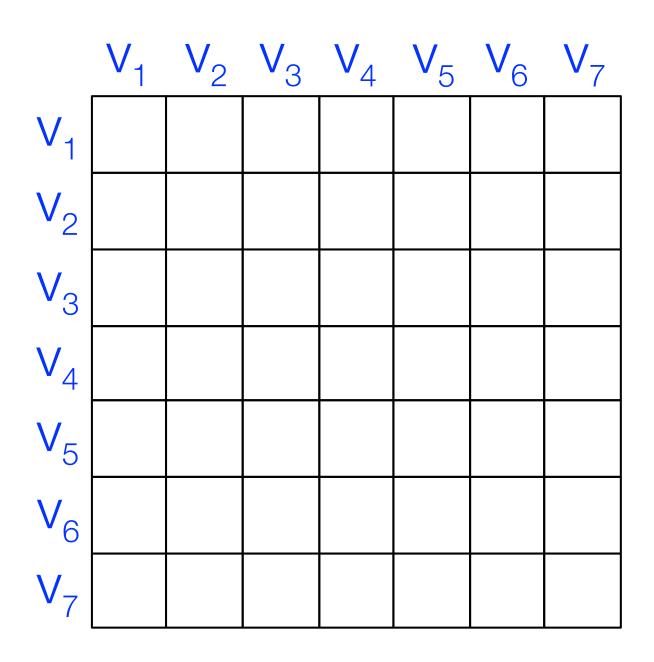




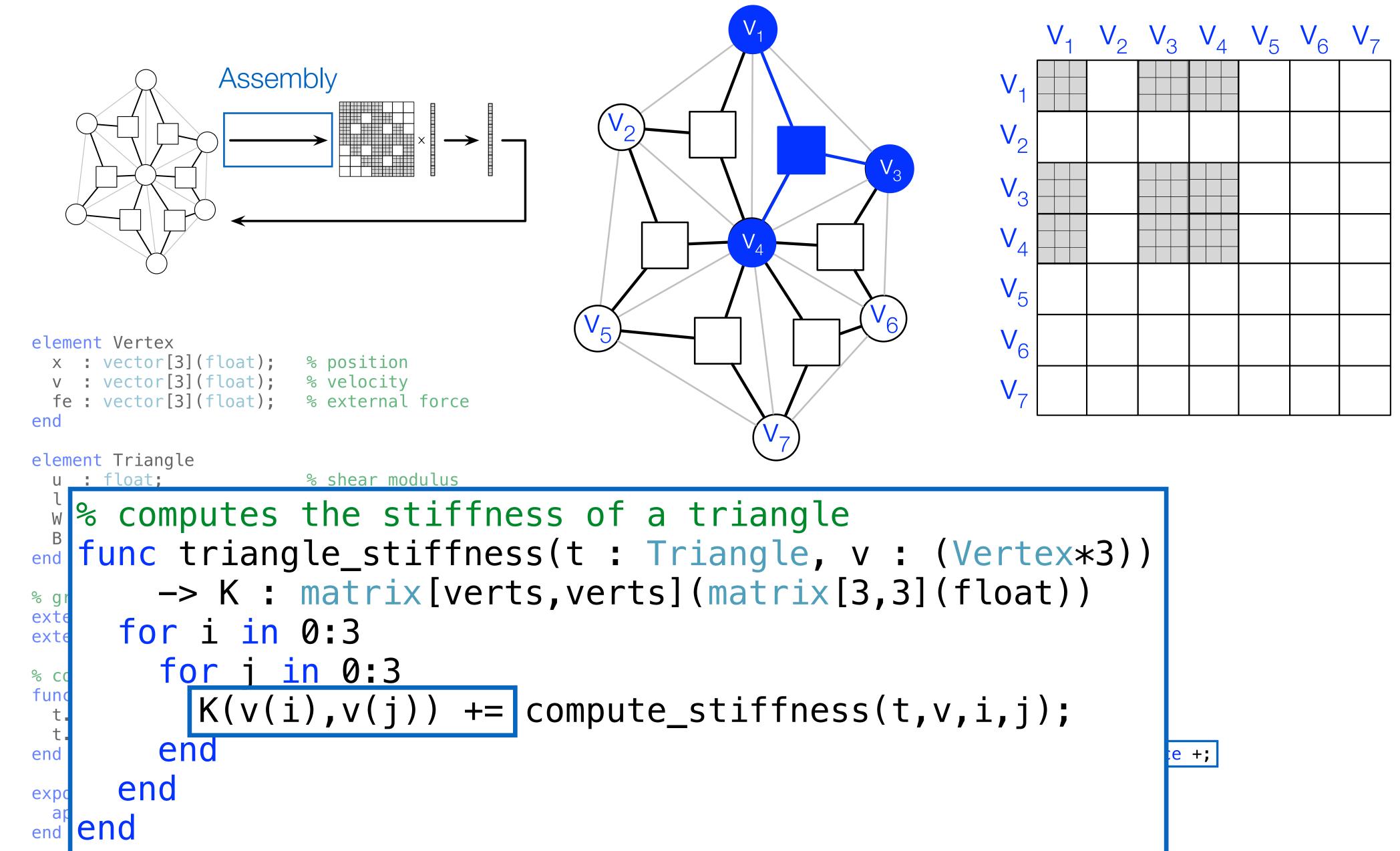
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element Vertex
 x : vector[3](float); % position
 v : vector[3](float); % velocity
 fe : vector[3](float); % external force
end
element Triangle
 u : float;
                      % shear modulus
                     % lame's first parameter
 l : float;
                        % volume
 W : float;
 B : matrix[3,3](float); % strain-displacement
end
% graph vertices and triangle hyperedges
            : set{Vertex};
extern verts
extern triangles : set{Triangle}(verts, verts, verts);
% compute triangle area
func compute_area(inout t : Triangle, v : (Vertex*3))
 t.B = compute_B(v);
 t.W = det(B) / 2.0;
end
export func init()
 apply compute_area to triangles;
```

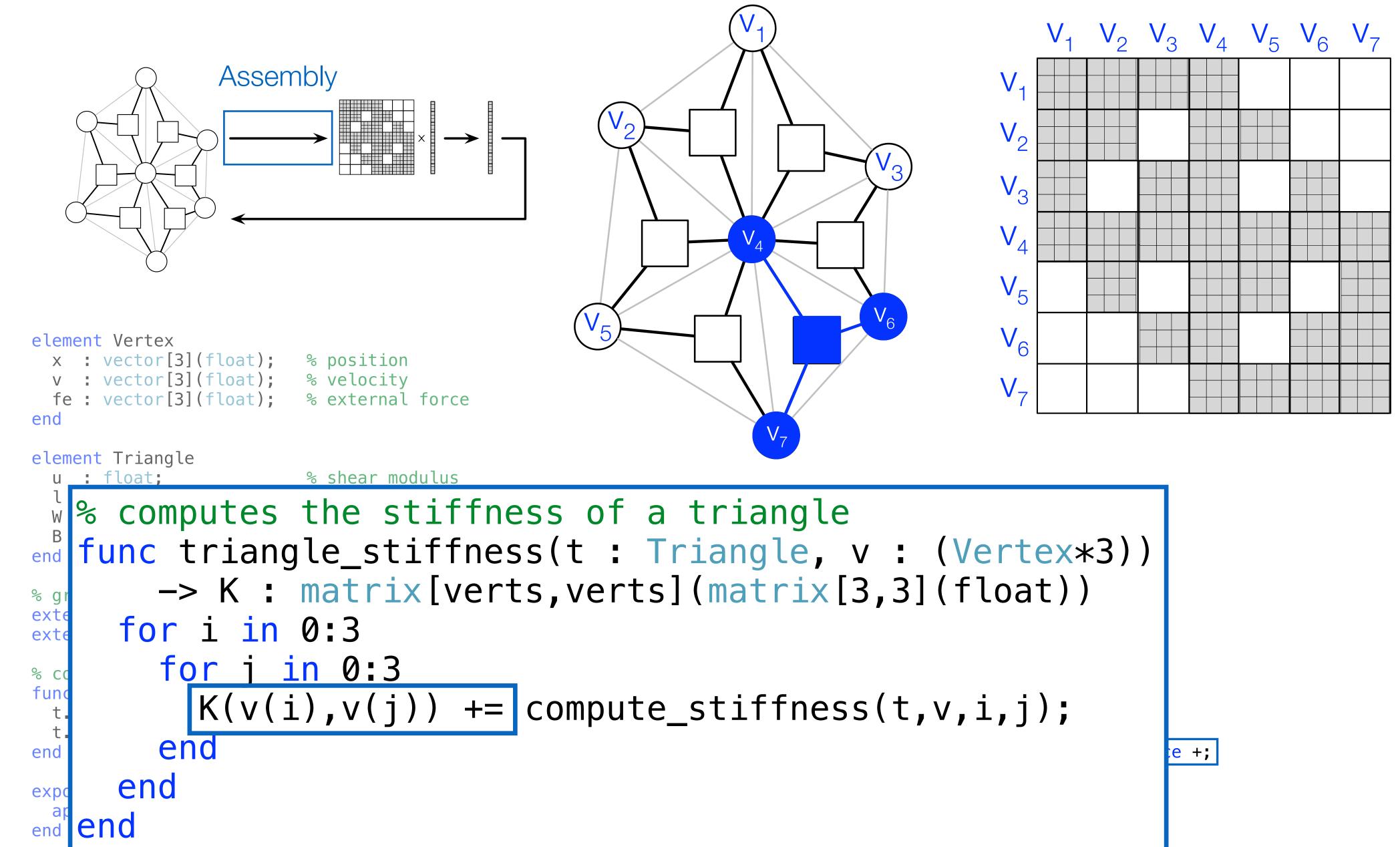
end

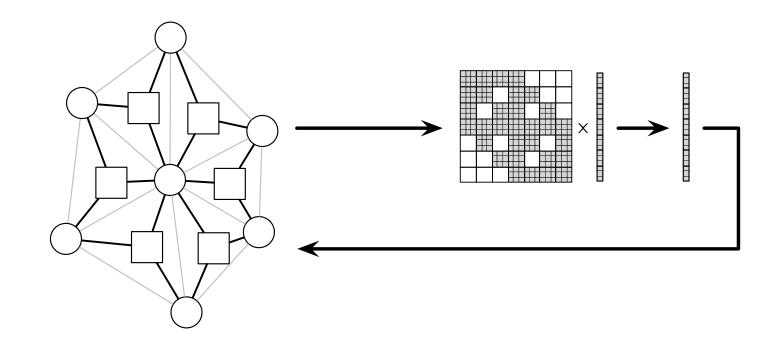




```
% newton's method
export func newton_method()
   while abs(f - verts.fe) > 1e-6
   K = map triangle_stiffness to triangles reduce +;
   // assemble force vector
   // compute new position
   end
end
```



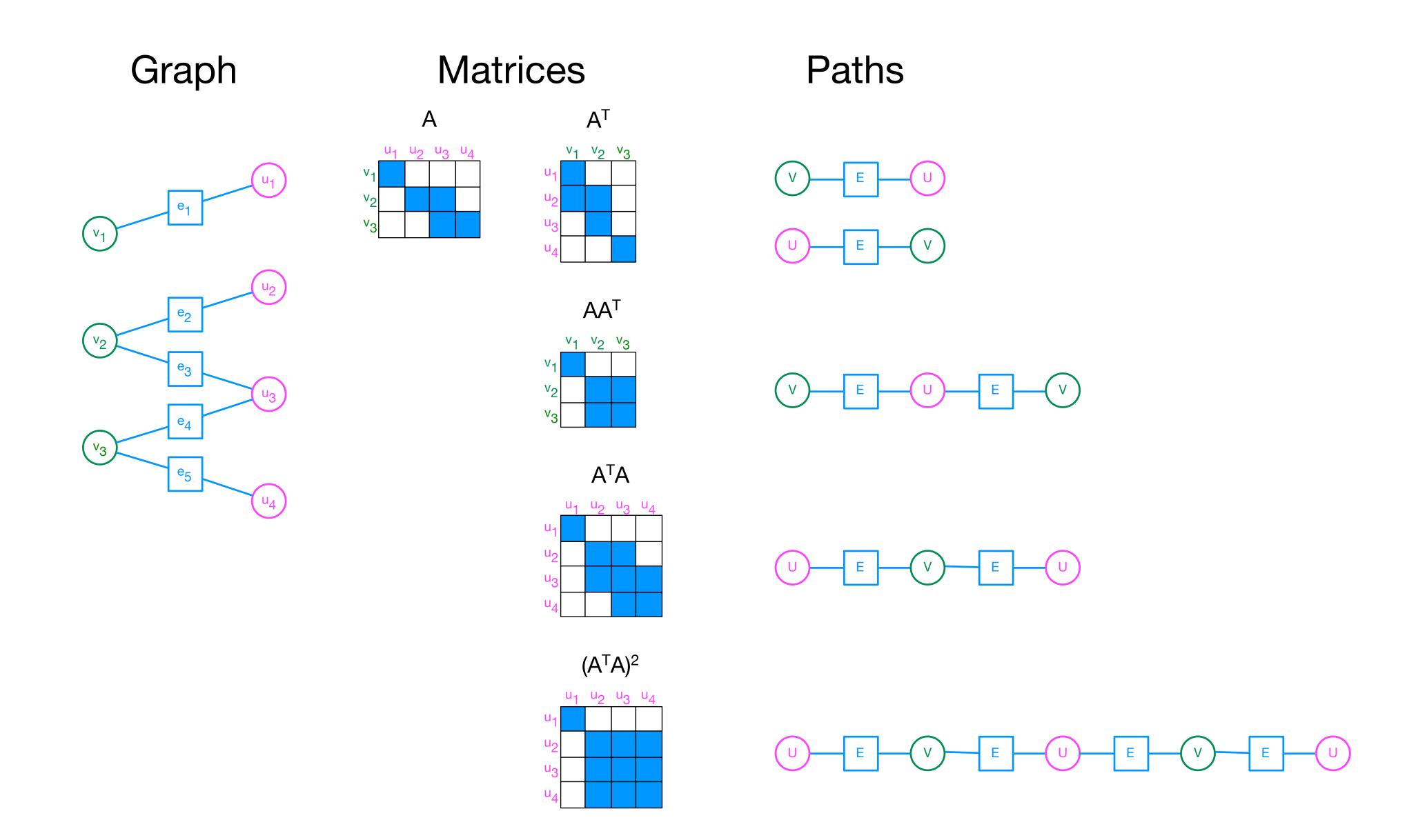




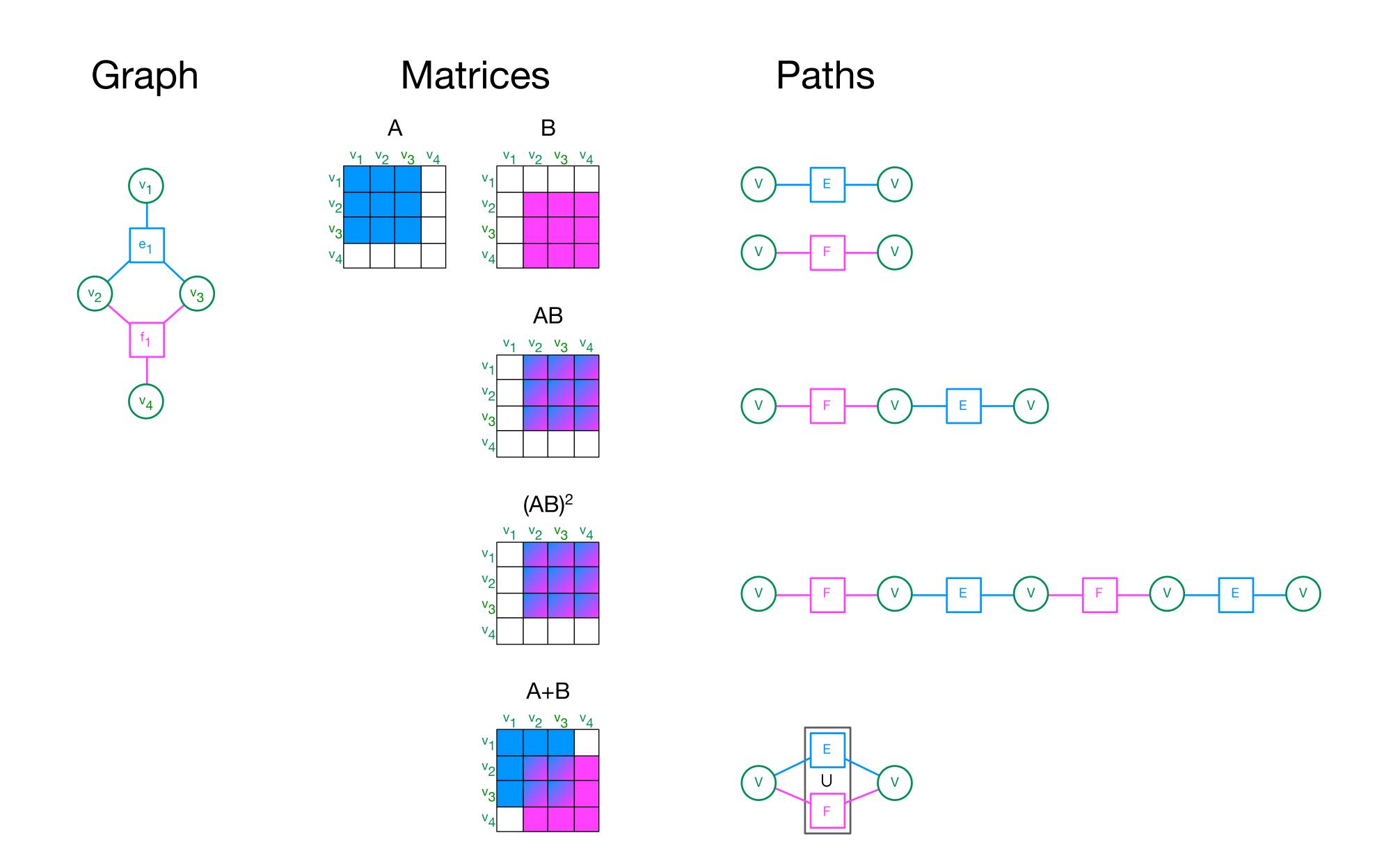
```
element Vertex
 x : vector[3](float);
                         % position
 v : vector[3](float);
                        % velocity
 fe : vector[3](float); % external force
end
element Triangle
 u : float;
                         % shear modulus
 l : float;
                         % lame's first parameter
                         % volume
 W : float;
    : matrix[3,3](float); % strain-displacement
end
% graph vertices and triangle hyperedges
                : set{Vertex};
extern verts
extern triangles : set{Triangle}(verts, verts);
% compute triangle area
func compute_area(inout t : Triangle, v : (Vertex*3))
 t_B = compute_B(v);
 t.W = det(B) / 2.0;
end
export func init()
  apply compute_area to triangles;
end
```

```
% computes the stiffness of a triangle
func triangle_stiffness(t : Triangle, v : (Vertex*3))
    -> K : matrix[verts, verts](matrix[3,3](float))
  for i in 0:3
    for j in 0:3
      K(v(i),v(j)) += compute_stiffness(t,v,i,j);
    end
  end
end
% computes the force of a triangle on its vertices
func triangle_force(t : Triangle, v : (Vertex*3))
    -> f : vector[verts](vector[3](float))
  for i in 0:3
    f(v(i)) += compute_force(t,v,i);
  end
end
% newton's method
export func newton_method()
  while abs(f - verts.fe) > 1e-6
    K = map triangle_stiffness to triangles reduce +;
   f = map triangle_force
                              to triangles reduce +;
    verts.x = verts.x + K \ (verts.fe - f);
 end
end
```

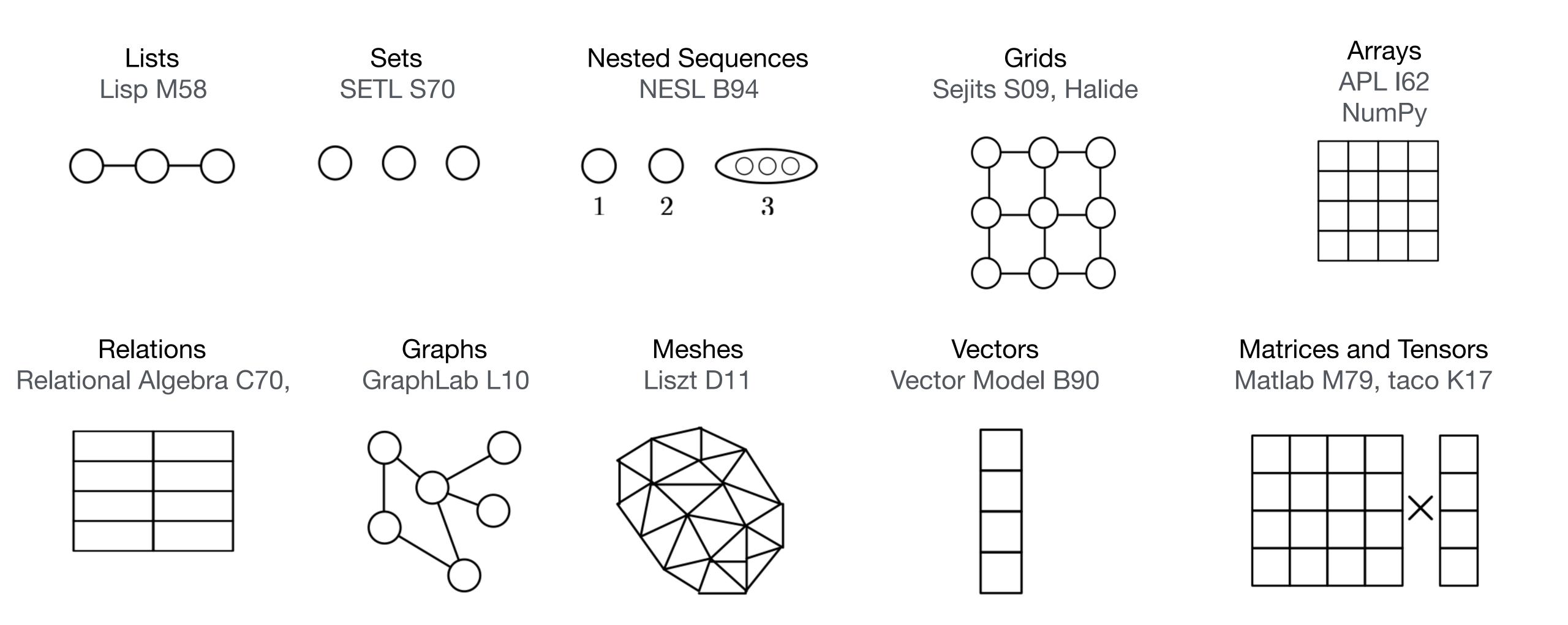
Paths through different sets



Paths through different hyperedge sets



Collection-Oriented Languages



A collection-oriented programming model provides collective operations on some collection/abstract data structure

Overview of lectures in the coming weeks

