Lecture 6 — Sparse Programming Systems

Stanford CS343D (Winter 2024)
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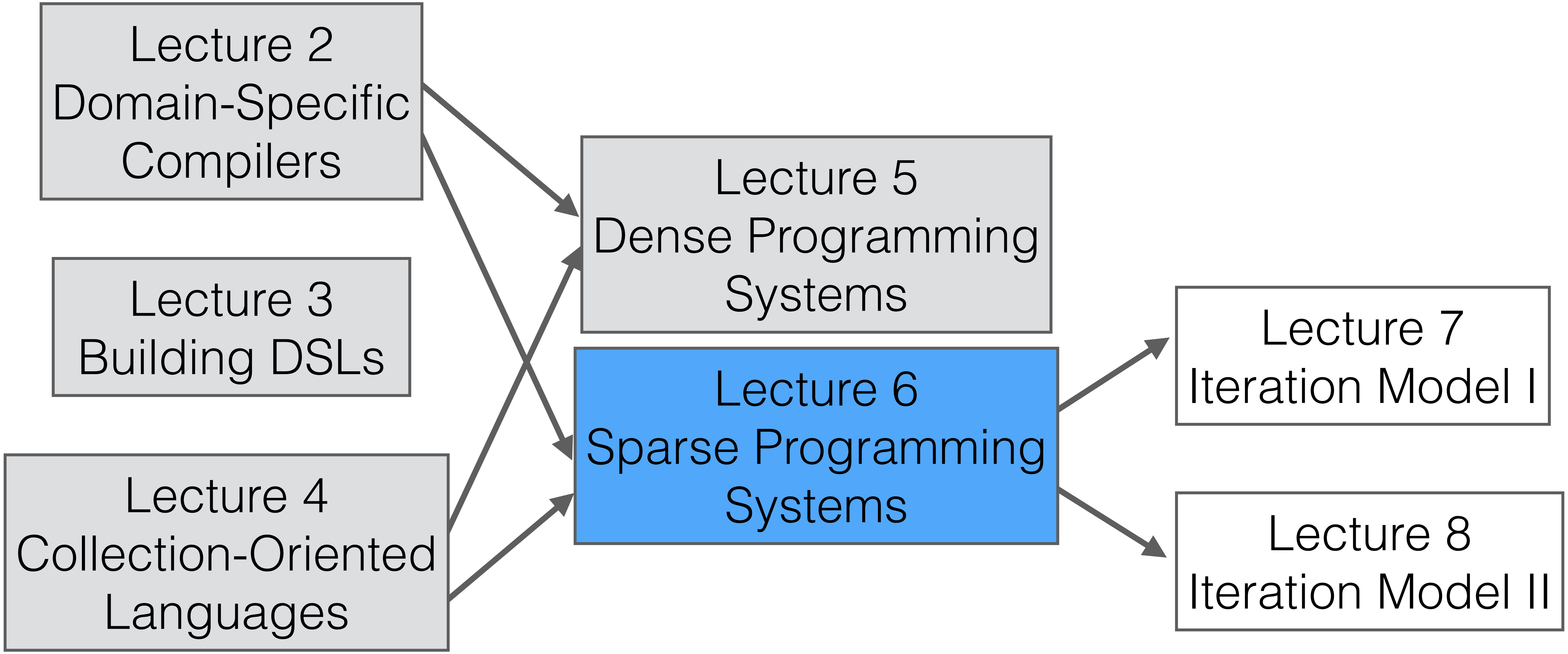
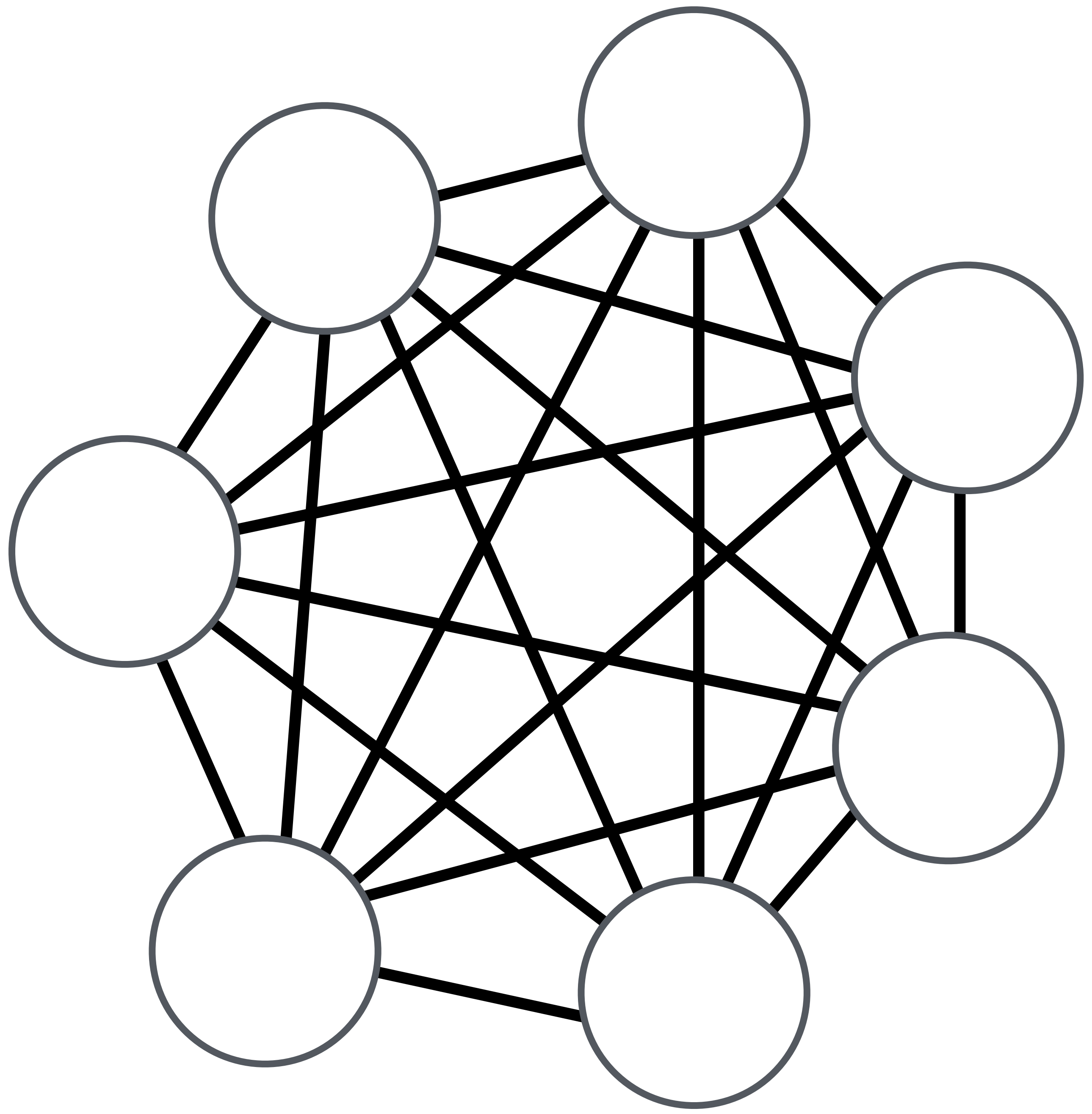
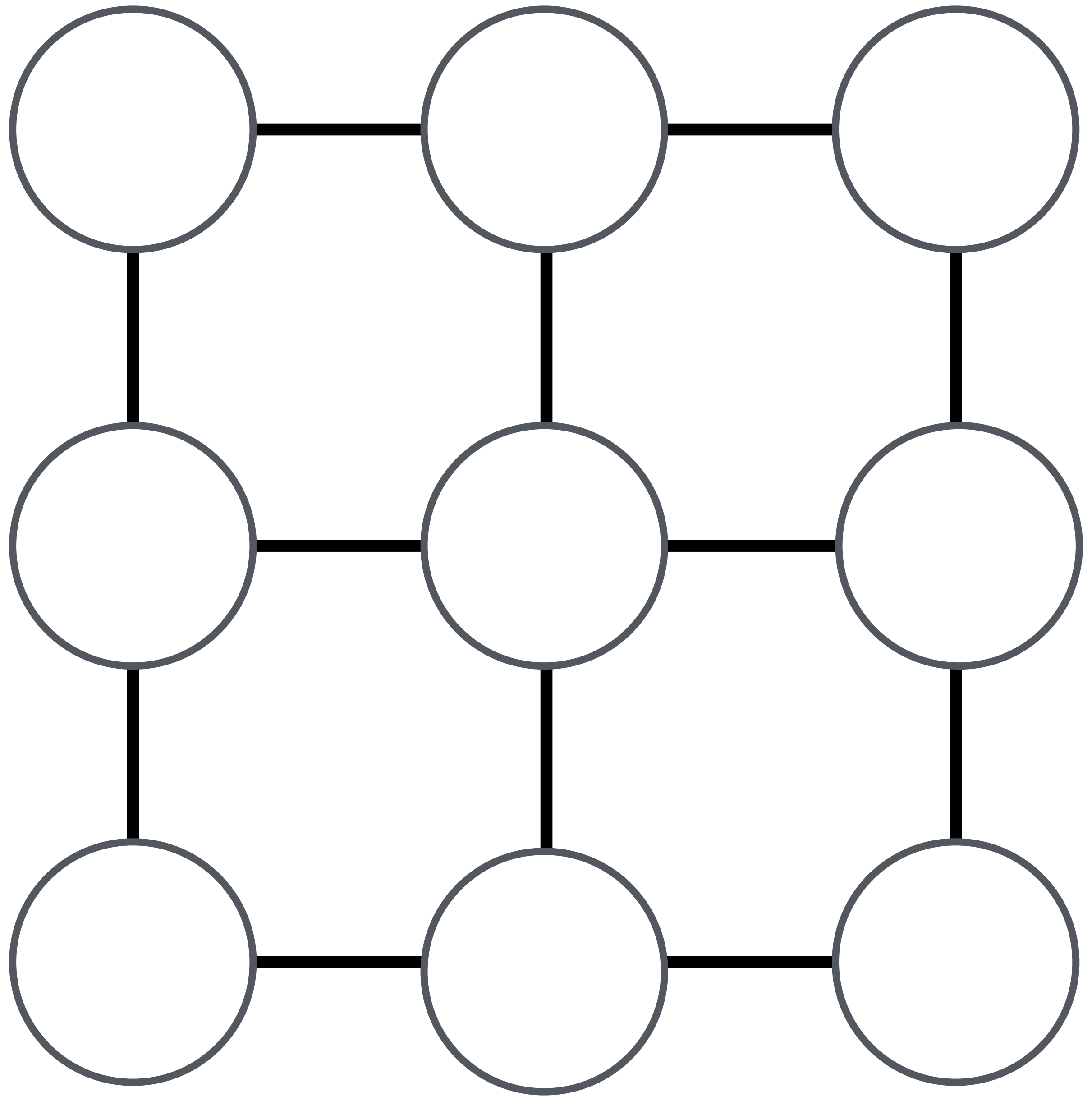

Lecture 2
Domain-Specific
Compilers
Lecture 3
Building DSLs
Lecture 4
Collection-Oriented
Languages
Lecture 5
Dense Programming
Systems
Lecture 6
Sparse Programming
Systems
Lecture 7
Iteration Model I
Lecture 8
Iteration Model II

Terminology: Regular and Irregular

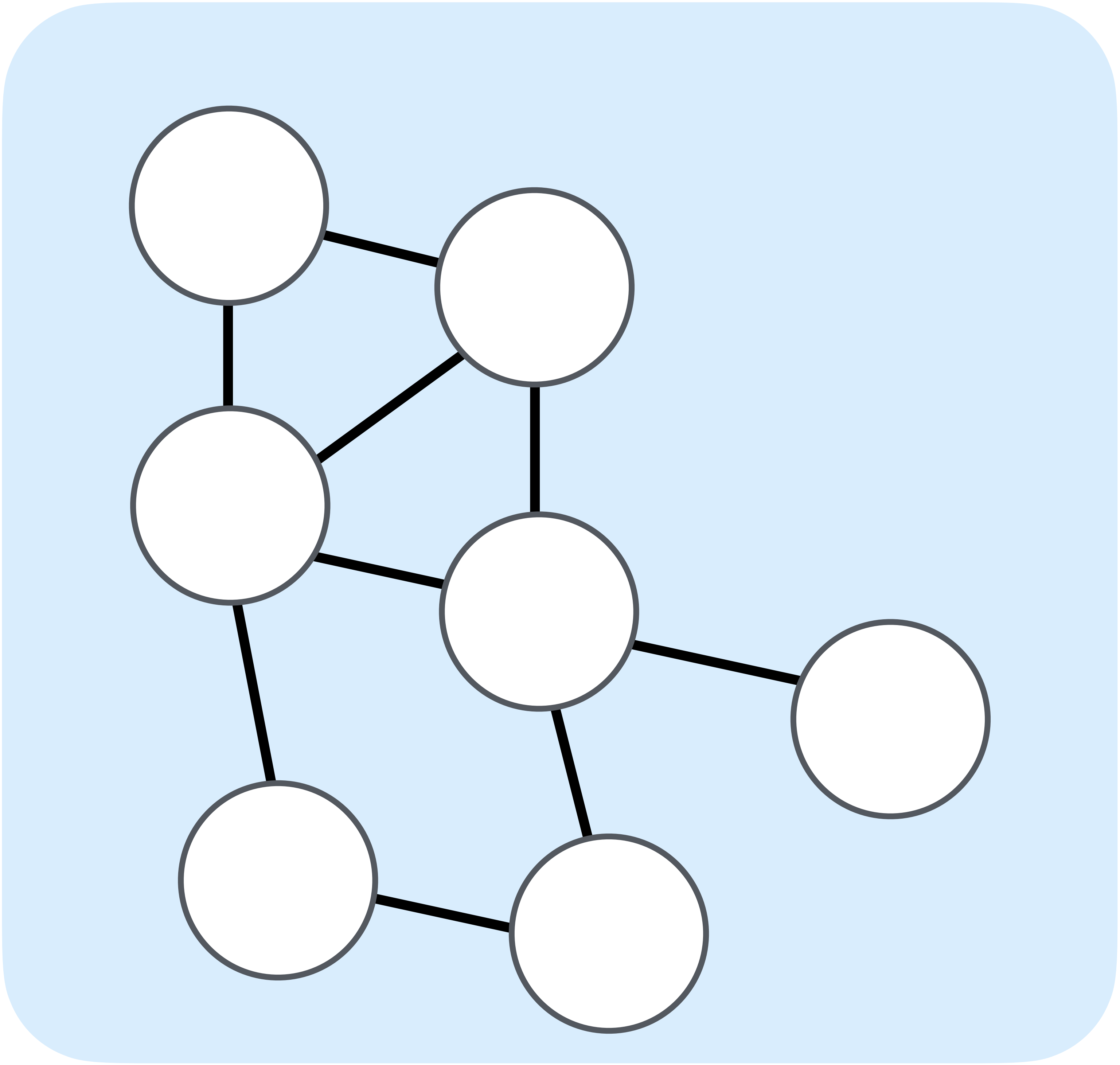
Fully Connected System



Regular System

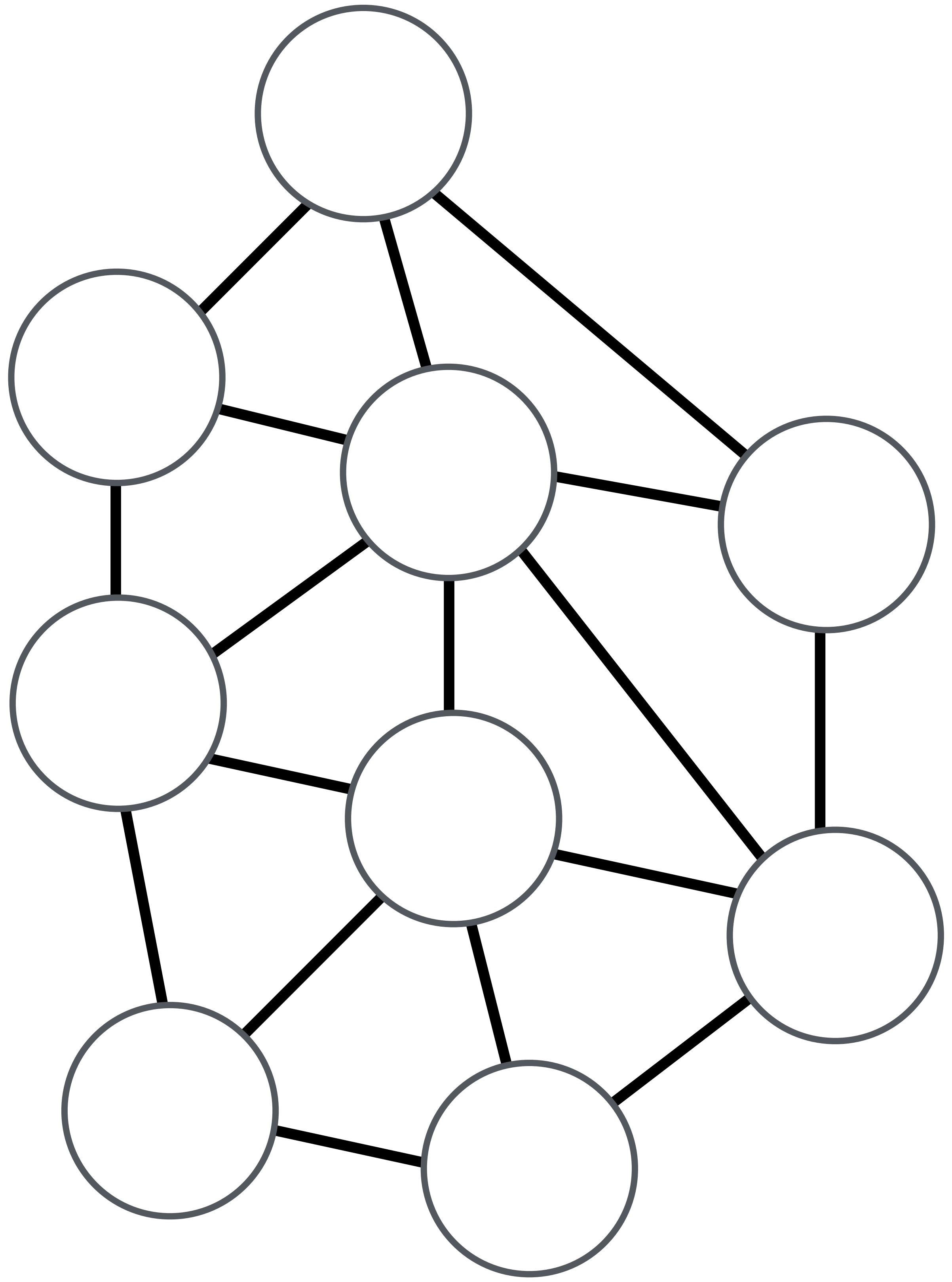


Irregular System

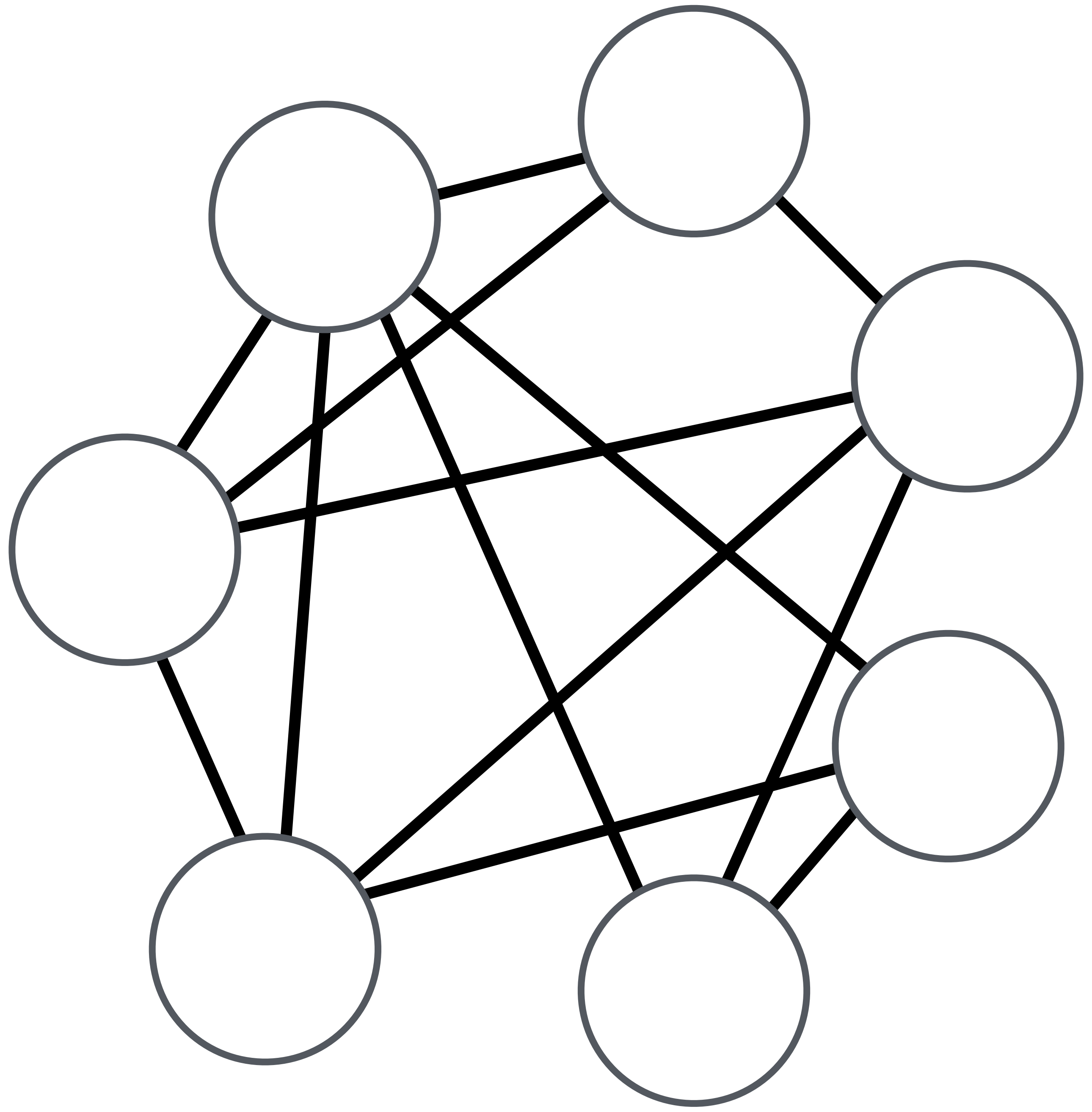


Three classes of irregular systems

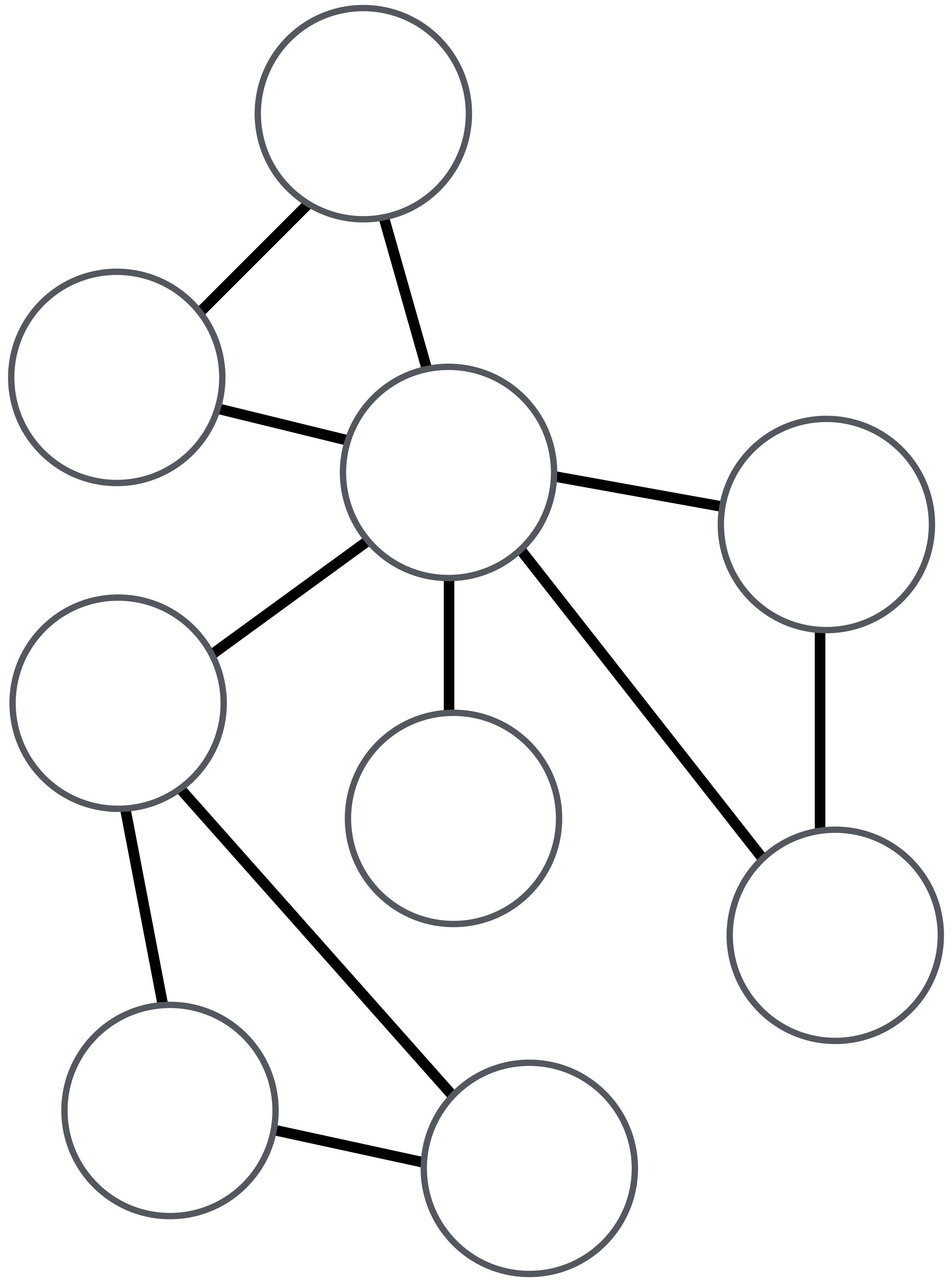
Road Networks



Fractional Sparsity

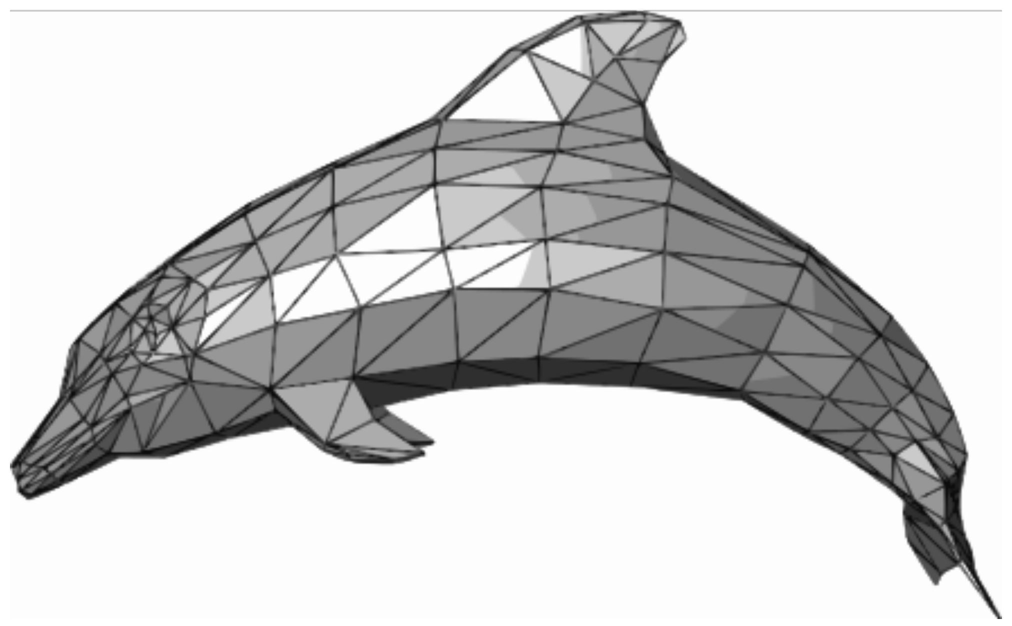


Power Law Graphs

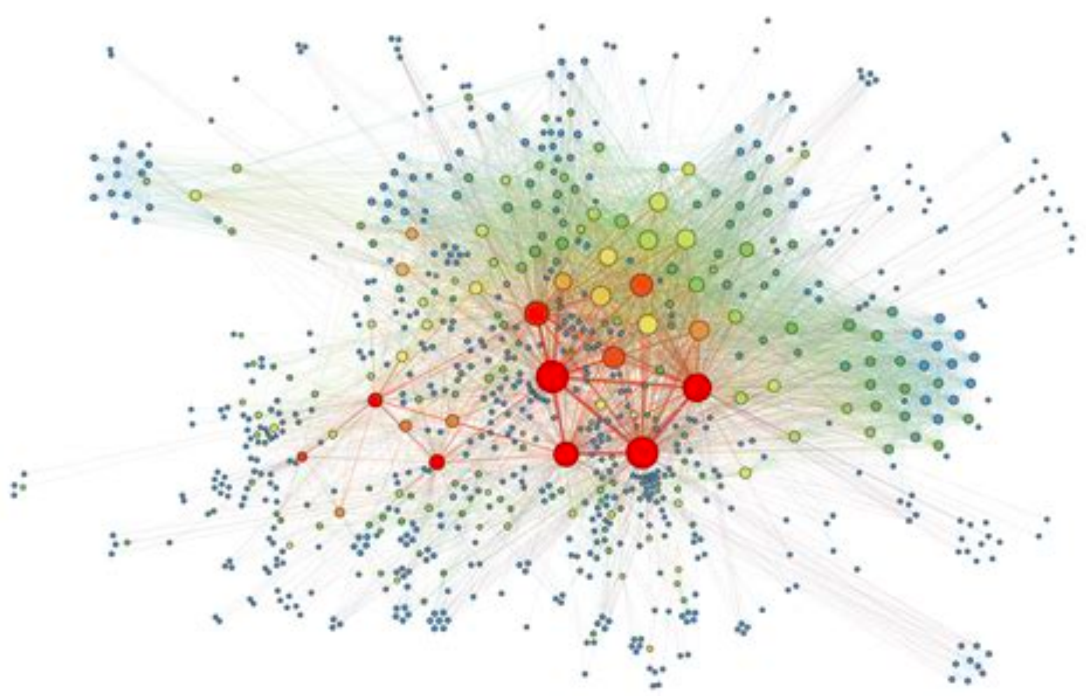


How sparse is graph/relational data? Often asymptotically sparse.

Conditioned Meshes



Power-law graphs



Assume an average degree of 150 (e.g., 150 friends)

Each matrix row then has 150 nonzeros

At 10,000 rows: $\frac{150 \cdot 10{,}000}{10{,}000^2} = 1.5\,\%$ nonzeros

At 100,000 rows: $\frac{150 \cdot 100{,}000}{100{,}000^2} = 0.15\,\%$ nonzeros

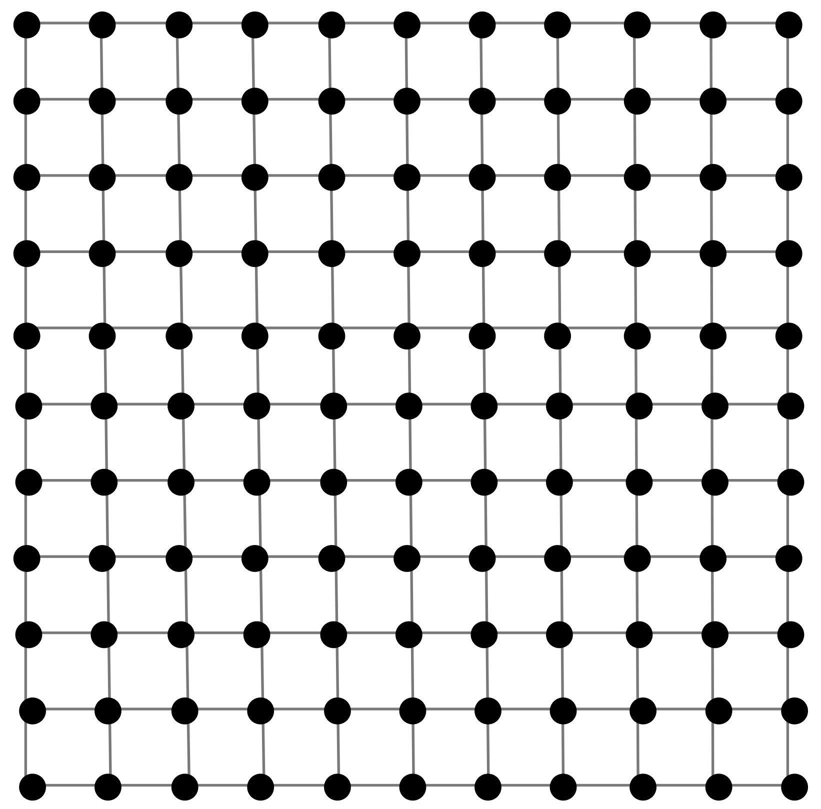
Matrix components: $O(n^2)$

Nonzeros: $O(n)$

Fraction of nonzeros: $O(1/n)$

Terminology: Dense and Sparse

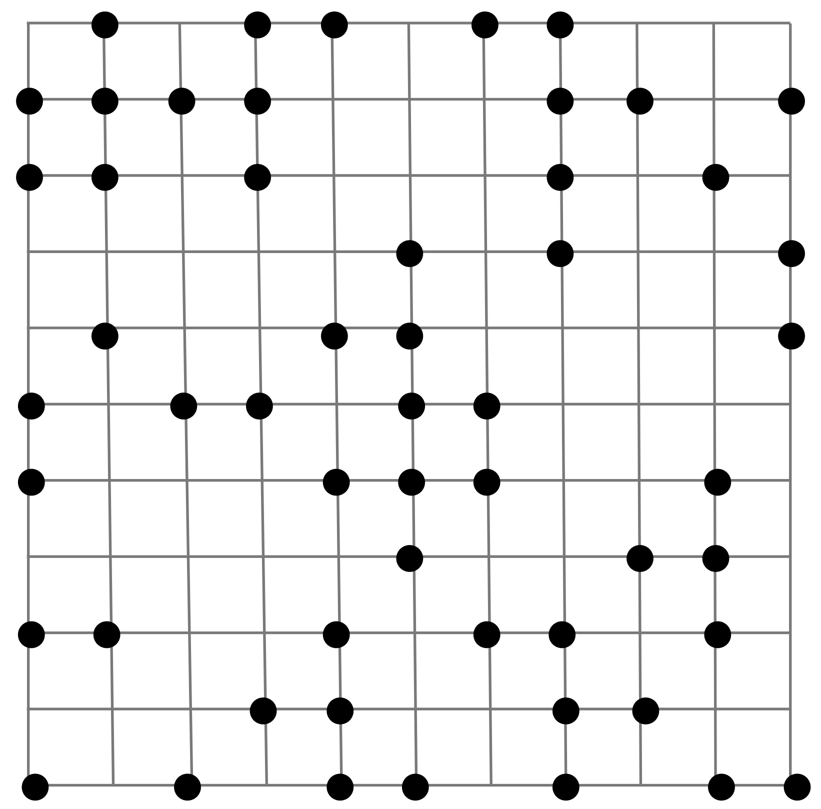
Dense loop iteration space



```
for (int i = 0; i < m; i++) {
  for (int j = 0; j < n; j++) {
    y[i] += A[i*n+j] * x[j];
  }
}
```

$$y = Ax$$

Sparse loop iteration space

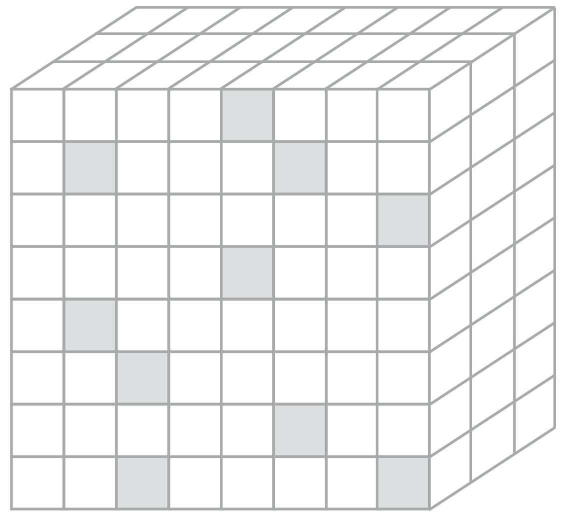


```
for (int i = 0; i < m; i++) {
 for (int pA = A2_pos[i]; pA < A_pos[i+1]; pA++) {
    int j = A_crd[pA];
    y[i] += A[pA] * x[j];
  }
}
```

$$y = Ax$$

Three sparse applications areas

Tensors



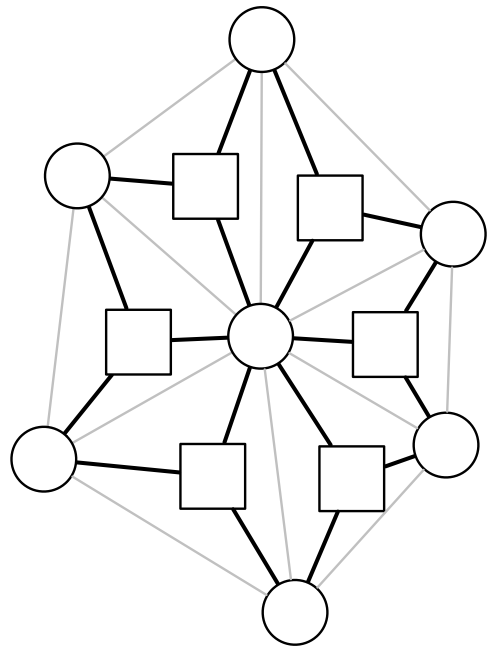
Nonzeros are a subset of the cartesian combination of sets

Relations

Names	City	Age
Peter	Boston	54
Mary	San Fransisco	35
Paul	New York	23
Adam	Seattle	84
Hilde	Boston	19
Bob	Chicago	76
Sam	Portland	32
Angela	Los Angeles	62

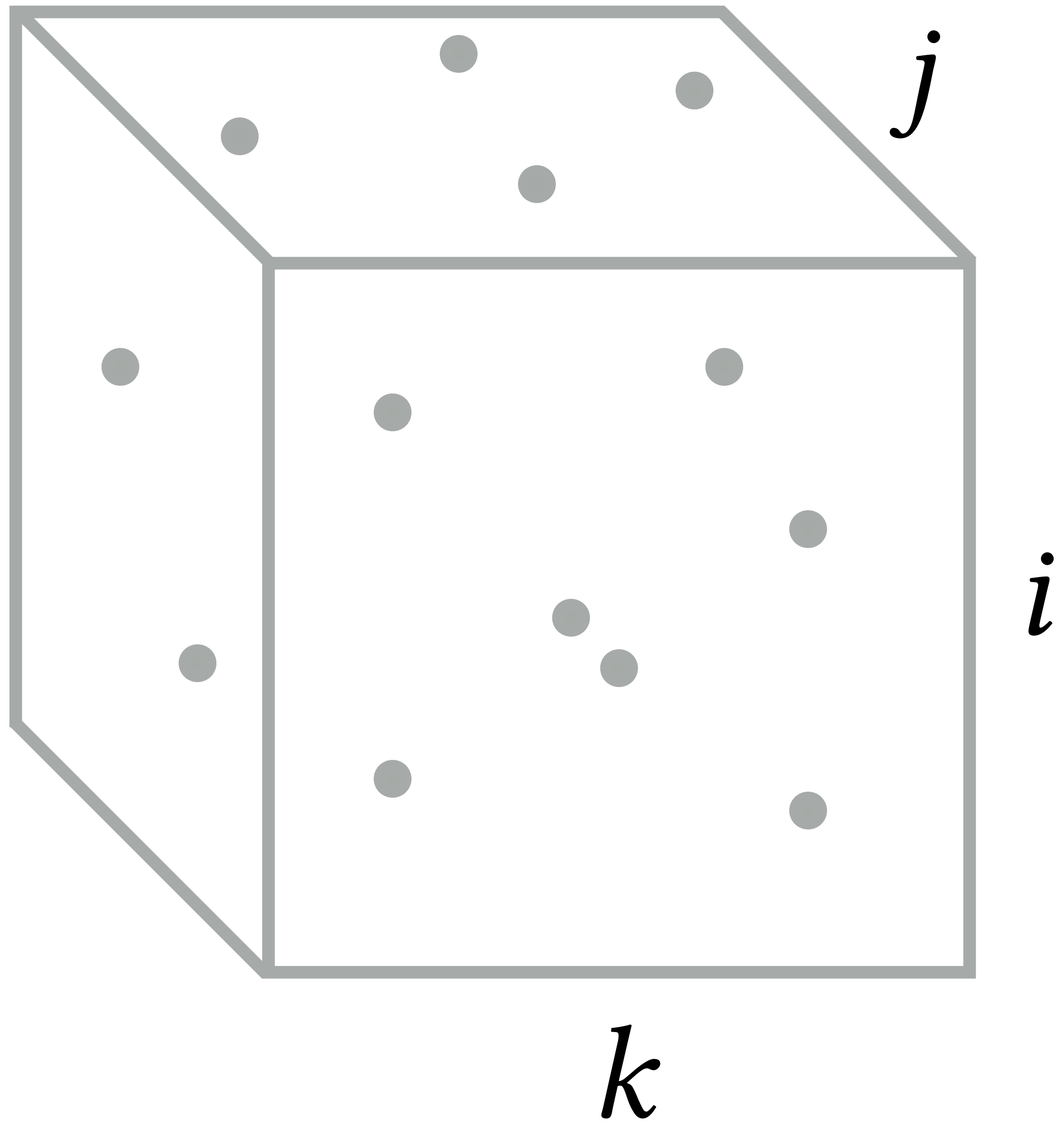
A relation is a subset of the cartesian combination of sets

Graphs



Graph edges are a subset of the cartesian combination of sets

Sparse Iteration Spaces



Relations, graphs, and tensors share a lot of structure but are specialized for different purposes

Relations

Names	City	Age
Peter	Boston	54
Mary	San Fransisco	35
Paul	New York	23
Adam	Seattle	84
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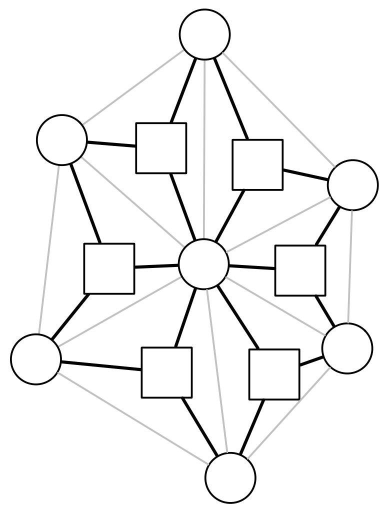
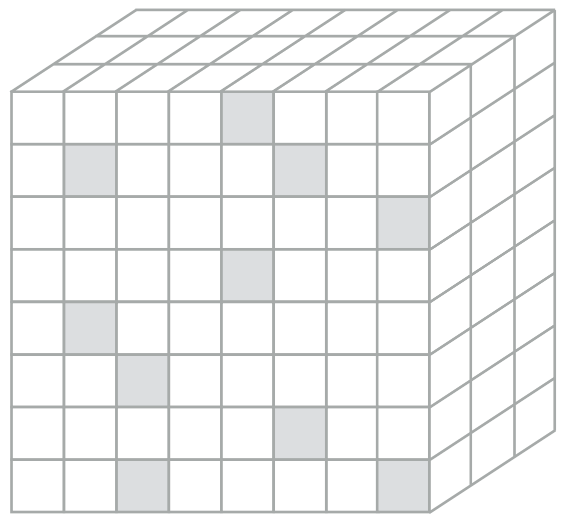
Combine data to form systems

Graphs

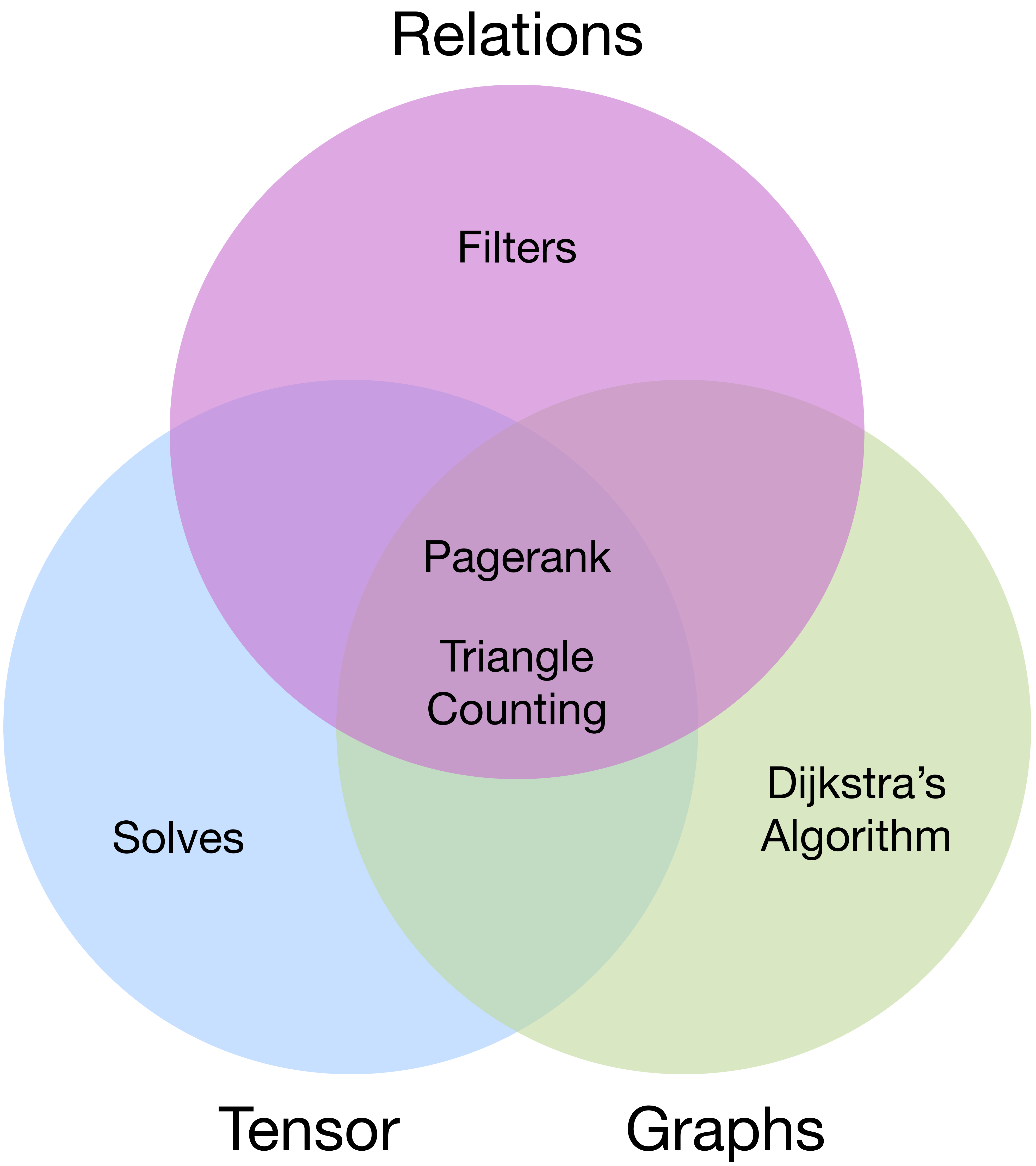


Local operations on systems

Tensors

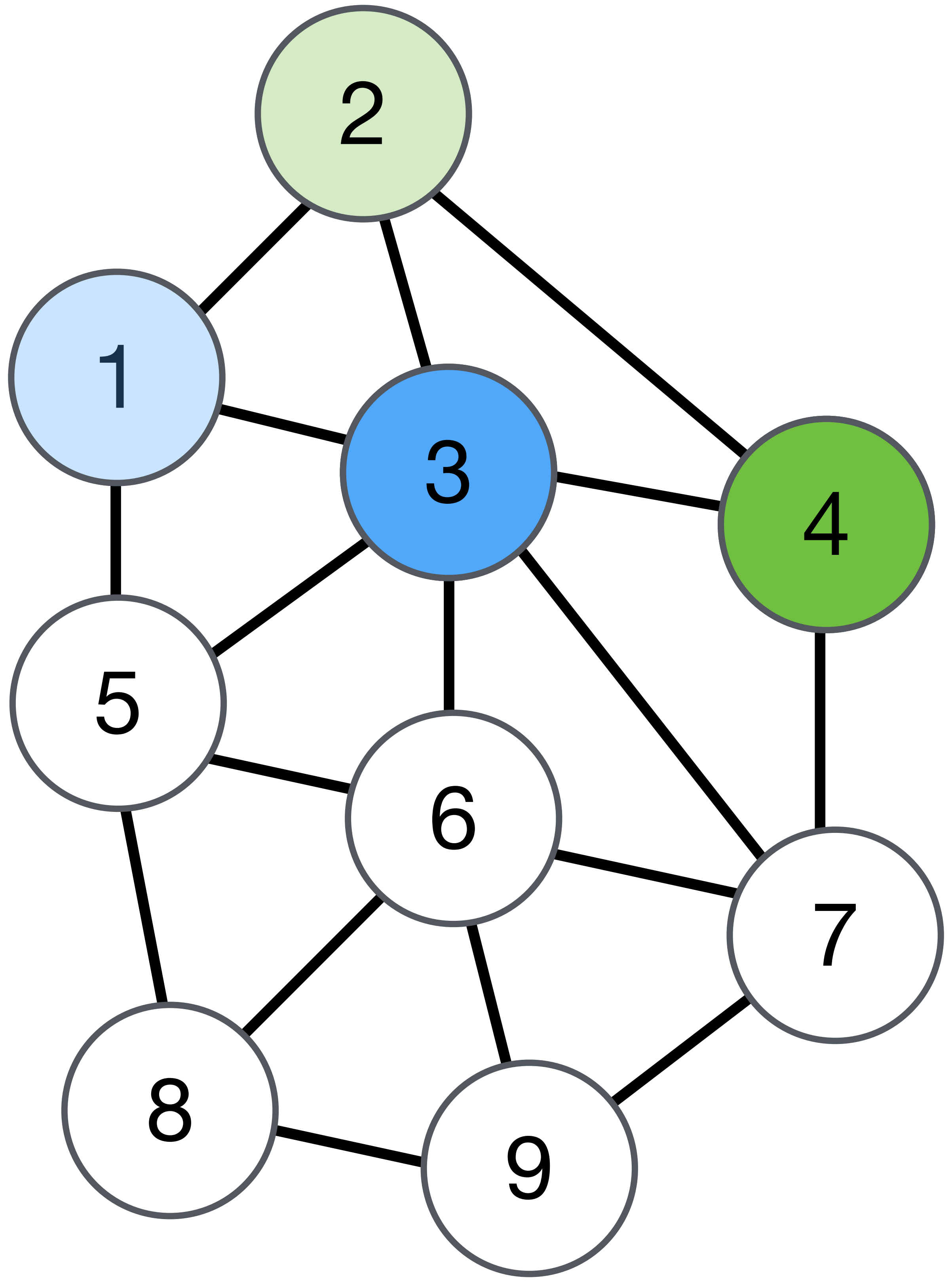


Global operations on systems



Triangle counting on graphs, relations, and tensors

On graphs



On relations

$$Q_{\triangle} = E(A, B) \bowtie E(B, C) \bowtie E(C, A)$$

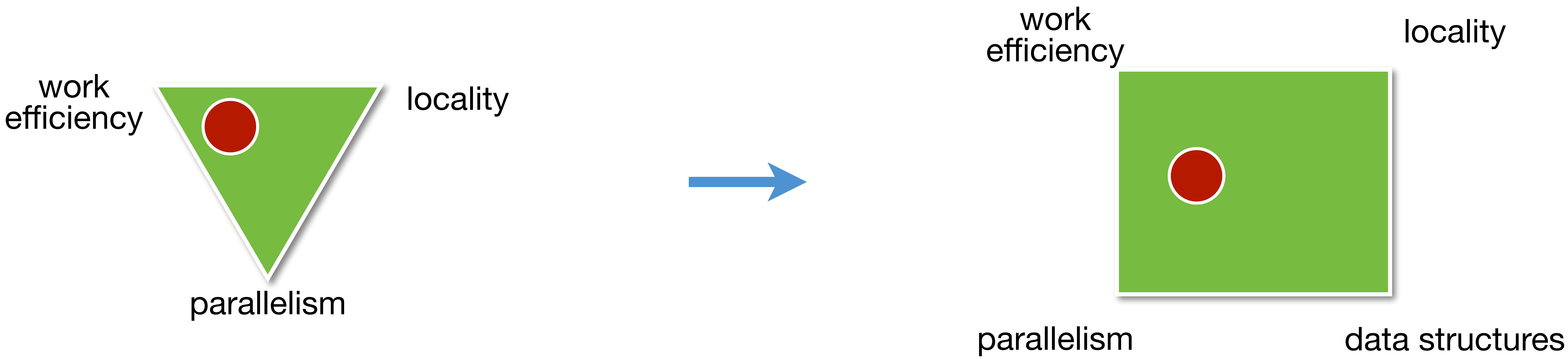
On tensors

$$\frac{1}{6}\text{trace}(A^3).$$

Some important developments in compilers and programming languages for sparse compilers

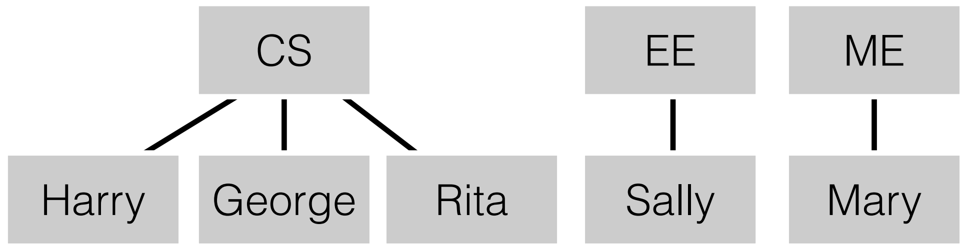
- 1960s: Development of libraries for sparse linear algebra
- 1970s: Relational algebra and the first relational database management systems: System R and INGRES
- 1980s: SQL is developed and has commercial success
- 1990s: Matlab gets sparse matrices and some dense to sparse linear algebra compilers are developed
- 2000s: Sparse linear algebra libraries for supercomputers and GPUs
- 2010s: Graph processing libraries become popular, compilers for databases, and compilers for sparse tensor algebra

Parallelism, locality, work efficiency still matters, but the key is choosing efficient data structures



Harry	CS
Sally	EE
George	CS
Mary	ME
Rita	CS

Harry	Sally	George	Mary	Rita
CS	EE	CS	ME	CS

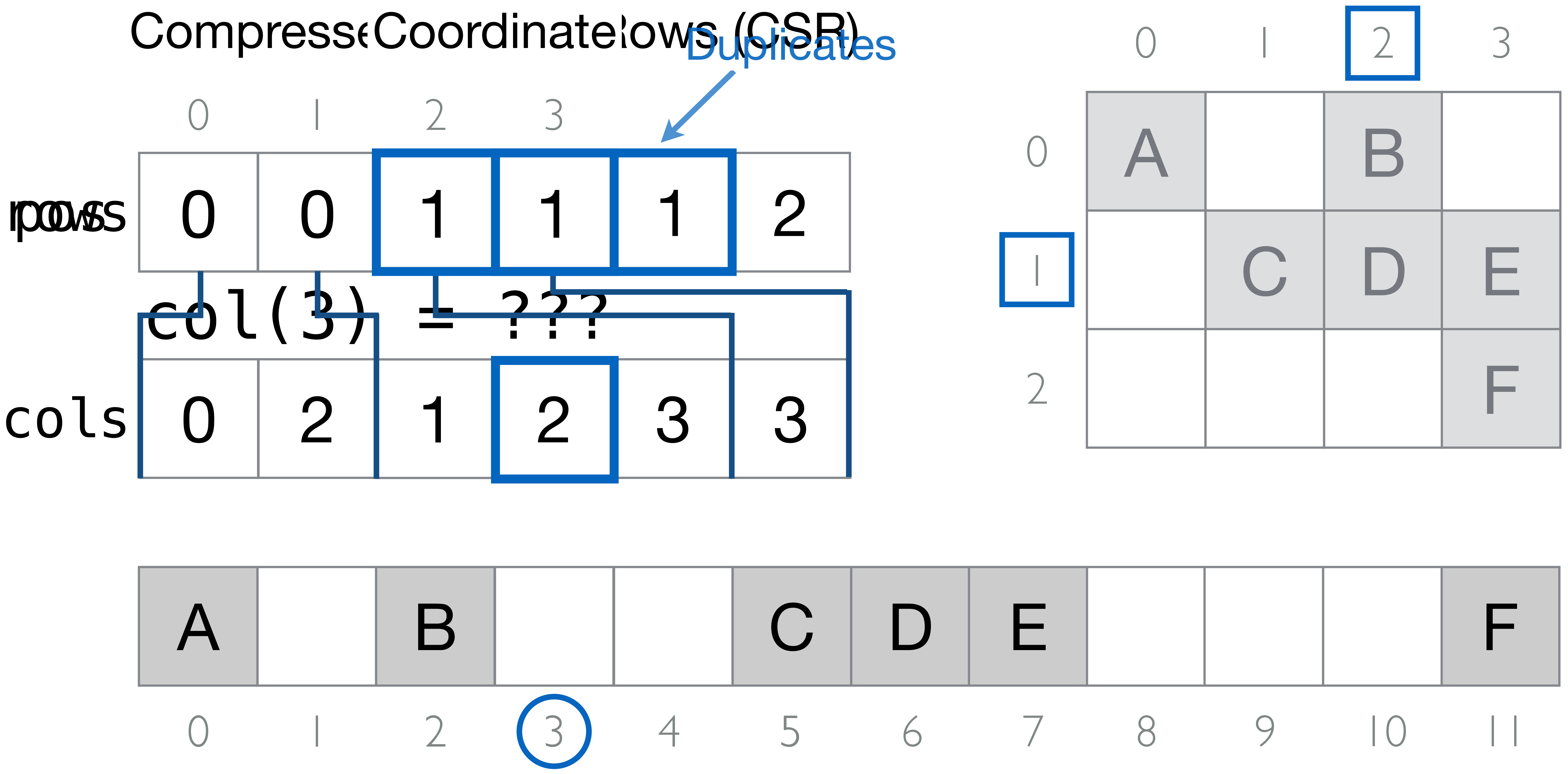


Sparse data structures in graphs, tensors, and relations encode coordinates in a sparse iteration space

Tensor (nonzeros)	Relation (rows)	Graph (edges)
(0,1)	(Harry,CS)	(v_1,v_5)
(2,3)	(Sally,EE)	(v_4,v_3)
(0,5)	(George,CS)	(v_5,v_3)
(5,5)	(Mary,ME)	(v_3,v_1)
(7,5)	(Rita,CS)	(v_3,v_5)

Values may be attached to these coordinates: e.g., nonzero values, edge attributes

Hierarchically compressed data structures (tries) reduce the number of values that need to be stored

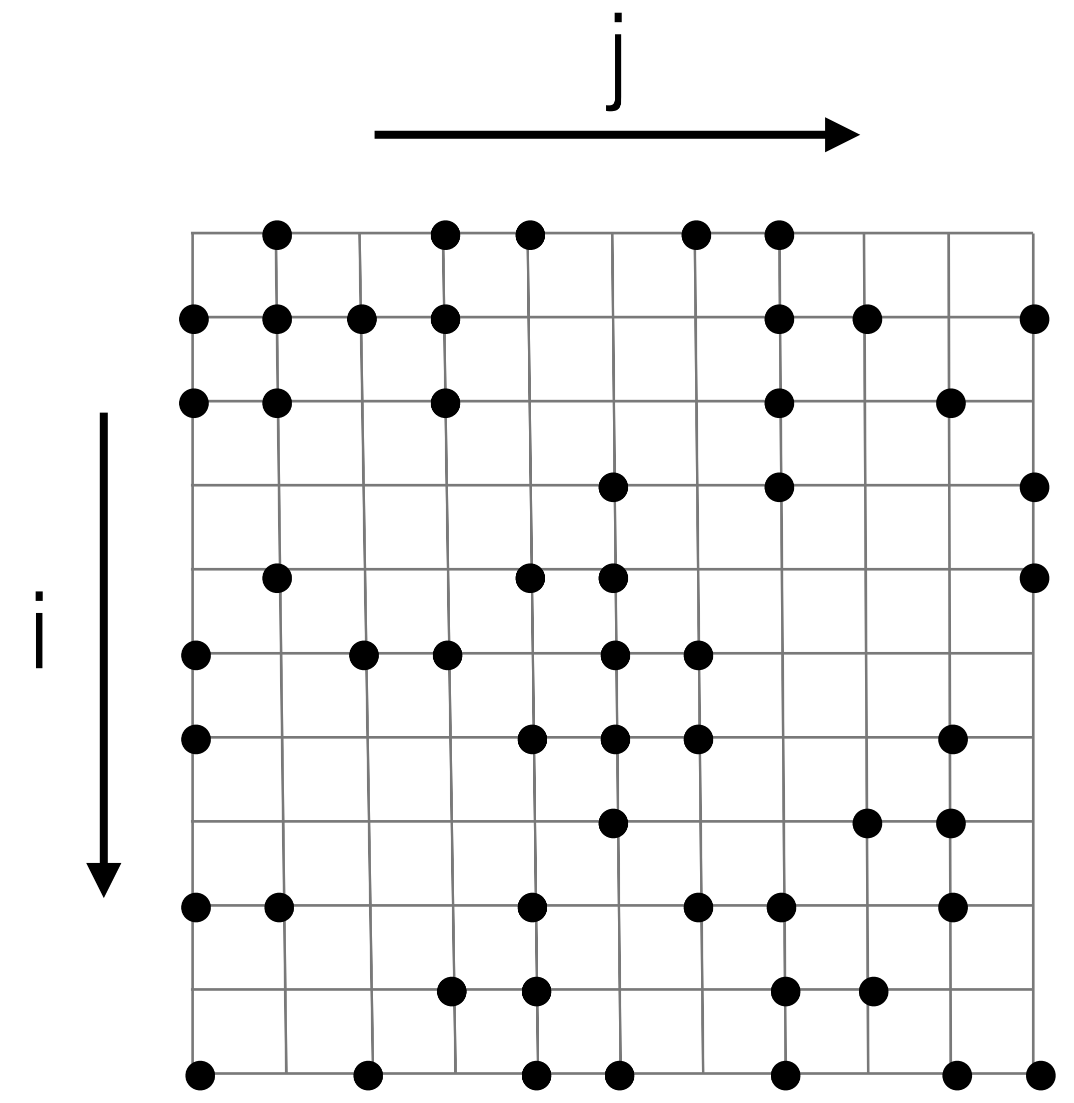


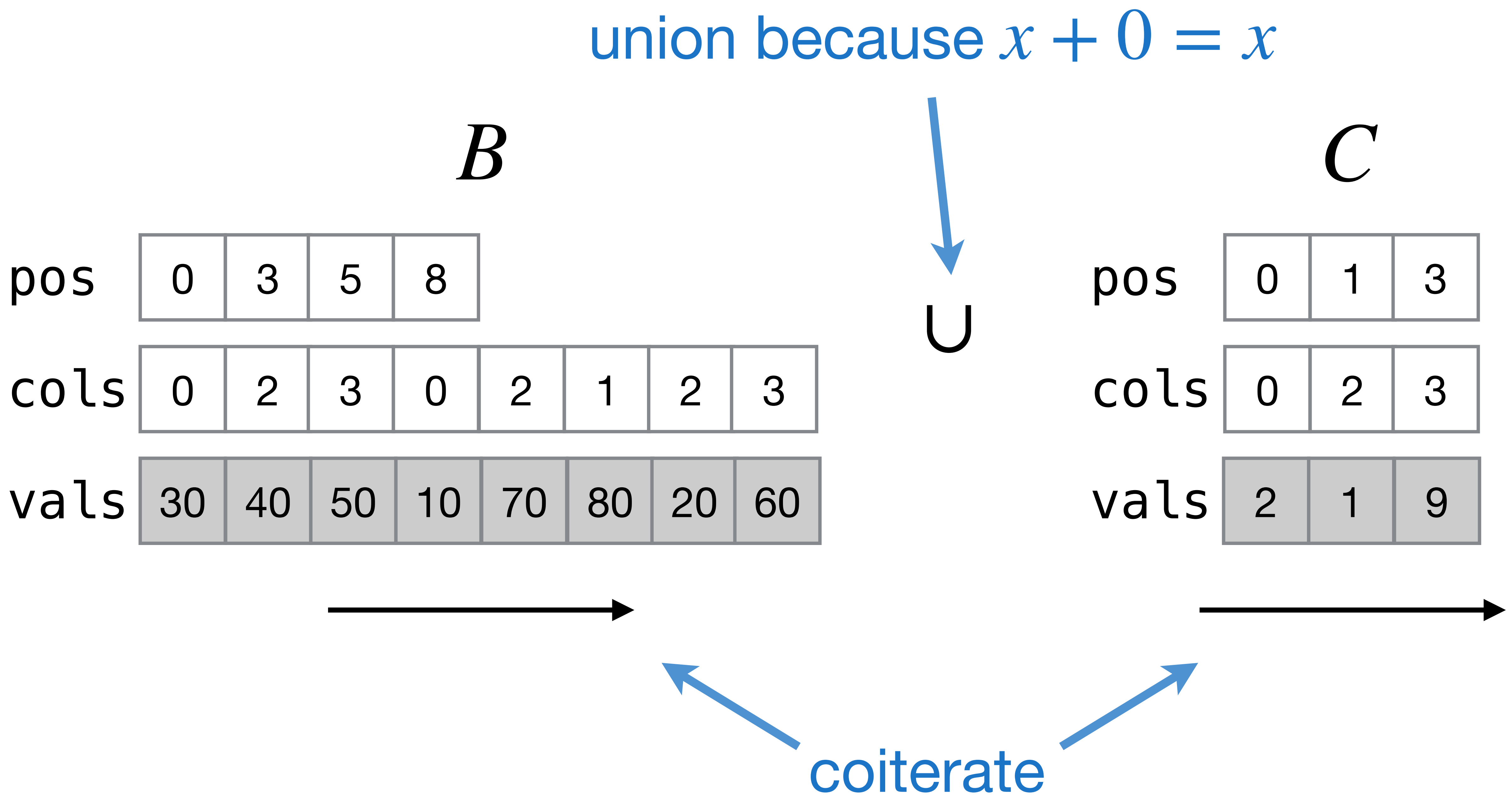
Iteration over sparse iteration spaces imply coiteration over sparse data structures

Linear Algebra: $A = B + C$

Tensor Index Notation: $A_{ij} = B_{ij} + C_{ij}$

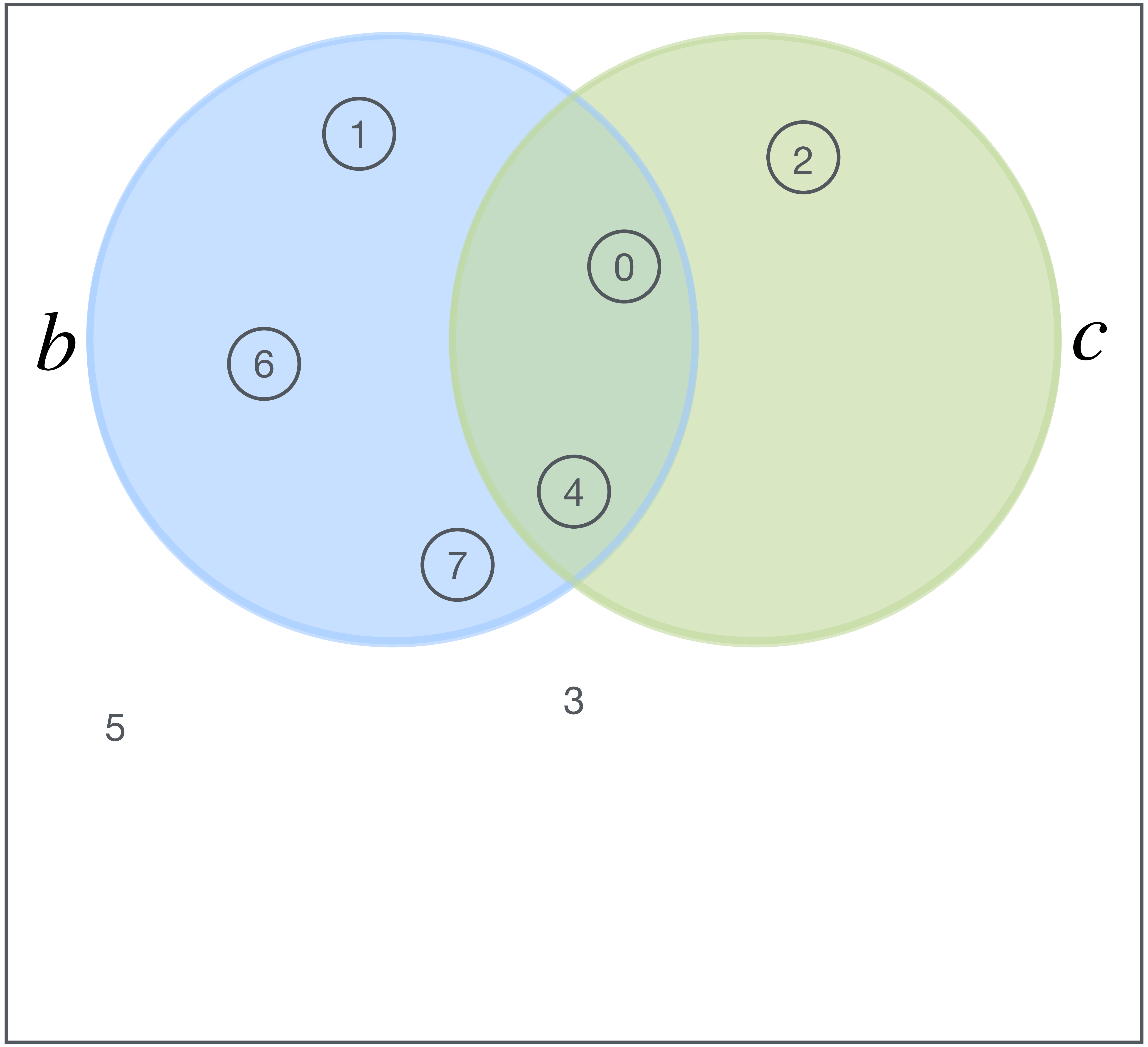
Iteration Space: $B_{ij} \cup C_{ij}$



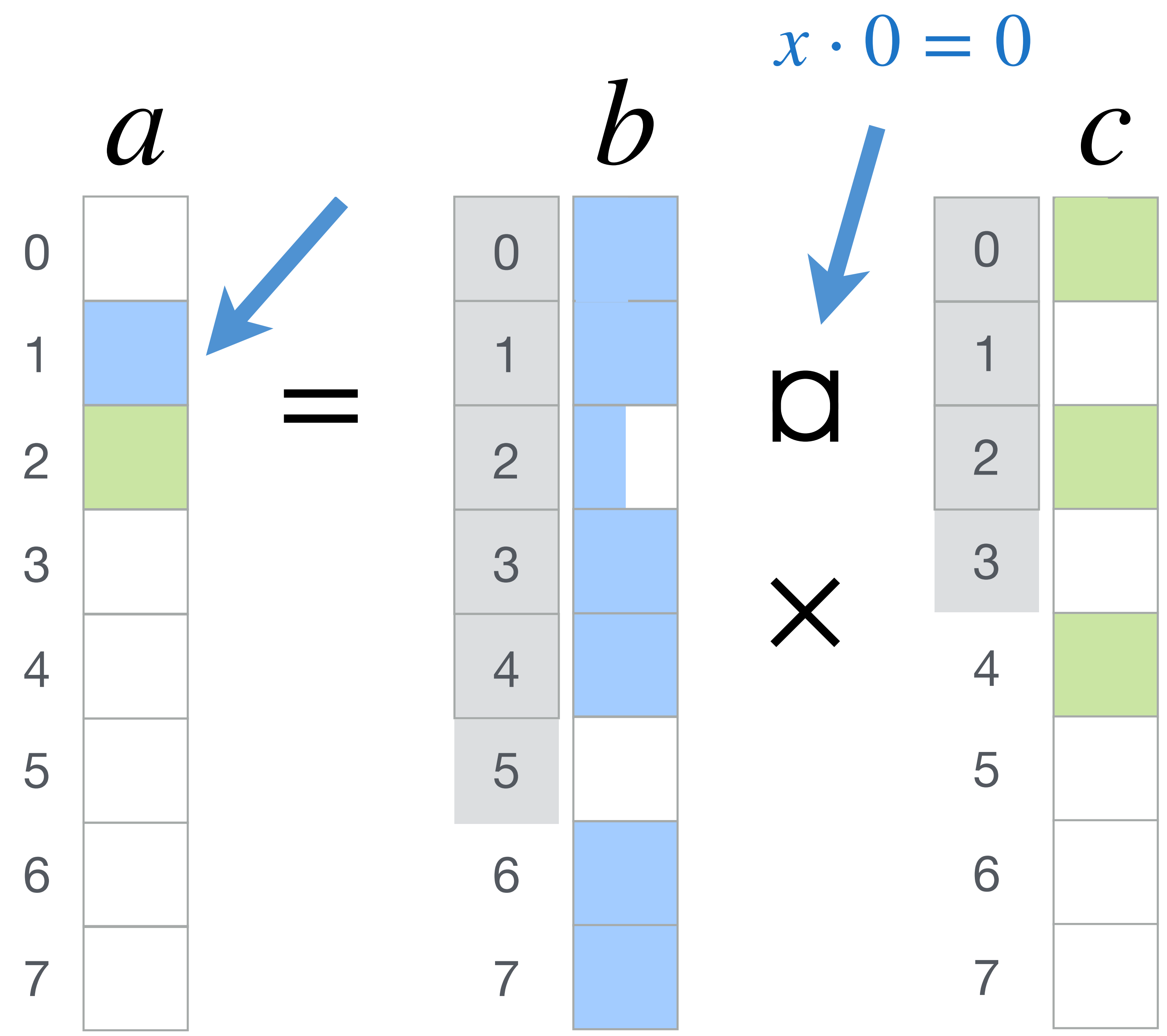


Merged coiteration

Coordinate Space



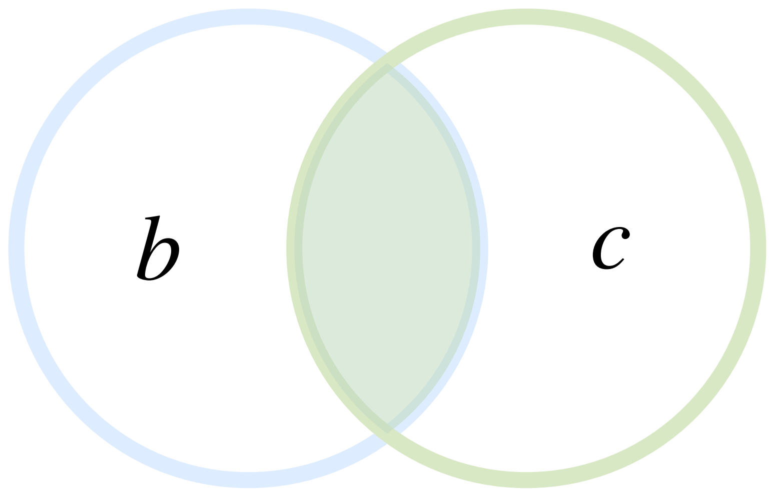
$$a_i = b_i + c_i$$



Merged coiteration code

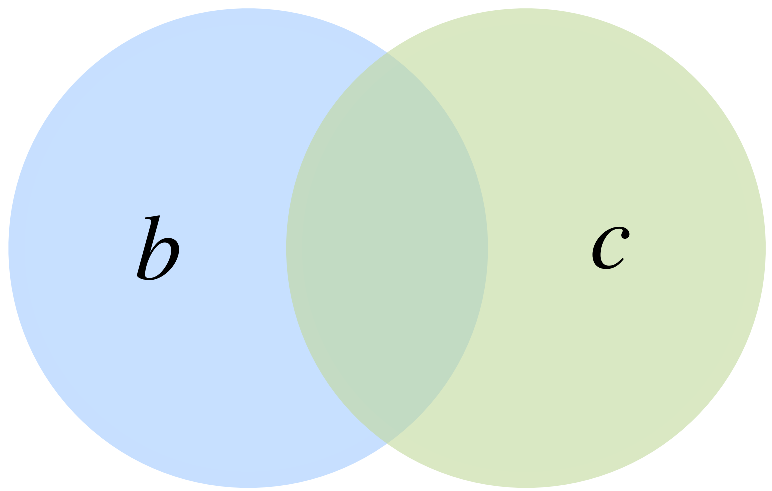
Intersection $b \cap c$

```
int pb = b_pos[0];
int pc = c_pos[0];
while (pb < b_pos[1] && pc < c_pos[1]) {
  int ib = b_crd[pb];
  int ic = c_crd[pc];
  int i = min(ib, ic);
  if (ib == i && ic == i) {
    a[i] = b[pb] * c[pc];
  }
  if (ib == i) pb++;
  if (ic == i) pc++;
}
```



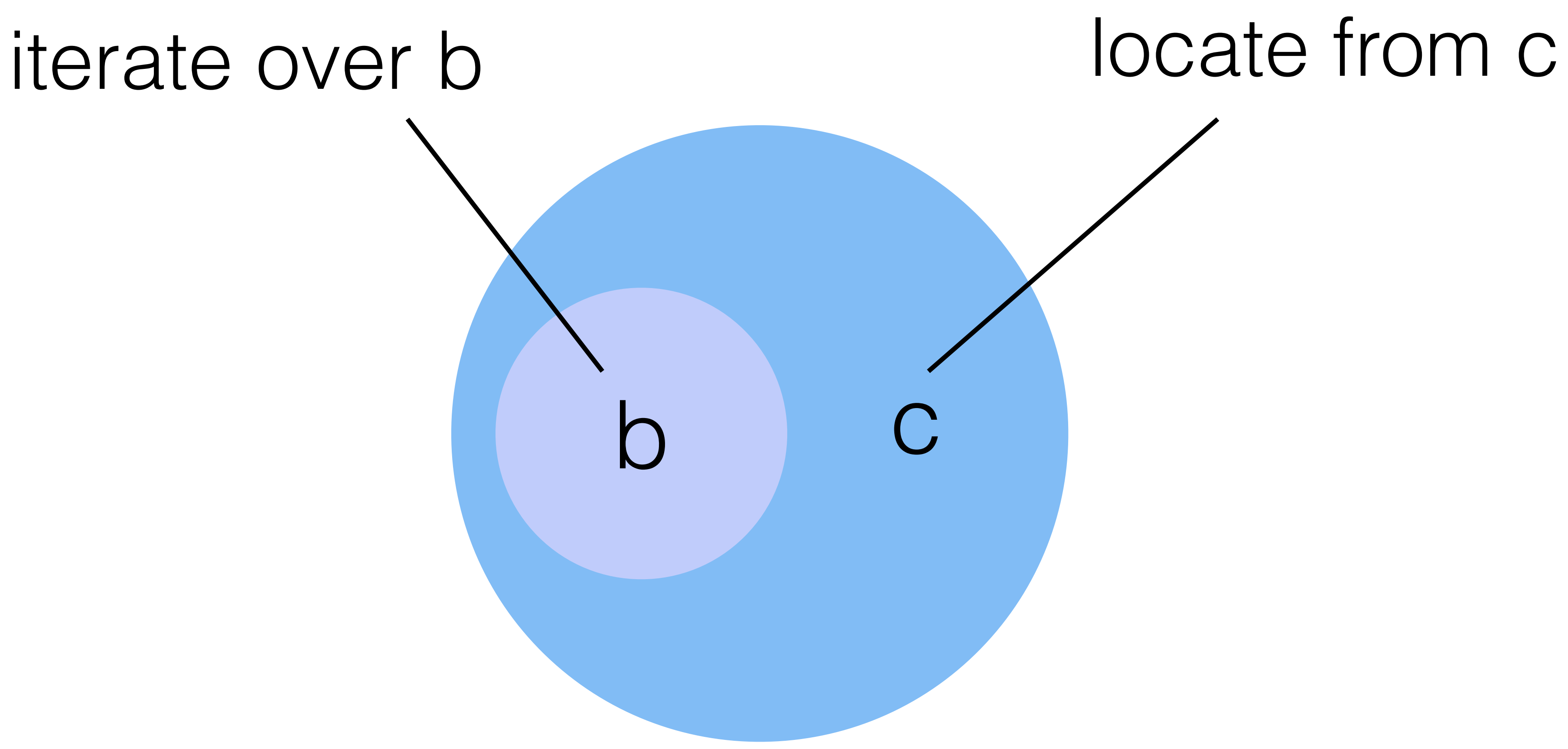
Union $b \cup c$

```
int pb = b_pos[0];
int pc = c_pos[0];
while (pb < b_pos[1] && pc < c_pos[1]) {
  int ib = b_crd[pb];
  int ic = c_crd[pc];
  int i = min(ib, ic);
  if (ib == i && ic == i) {
    a[i] = b[pb] + c[pc];
  }
  else if (ib == i) {
    a[i] = b[pb];
  }
  else {
    a[i] = c[pc];
  }
  if (ib == i) pb++;
  if (ic == i) pc++;
}

while (pb < b_pos[1]) {
  int i = b_crd[pb];
  a[i] = b[pb++];
}

while (pc < c_pos[1]) {
  int i = c_crd[pc];
  a[i] = c[pc++];
}
```



Iterate-and-locate examples (intersection)

$$a = \sum_i b_i c_i$$




```
for (int pb = b_pos[0]; pb < b_pos[1]; pb ++) {
  int i = b_crd[pb];
  a += b[pb] * c[i];
}
```

Separation of Algorithm, Data Representation, and Schedule

Algorithm Language

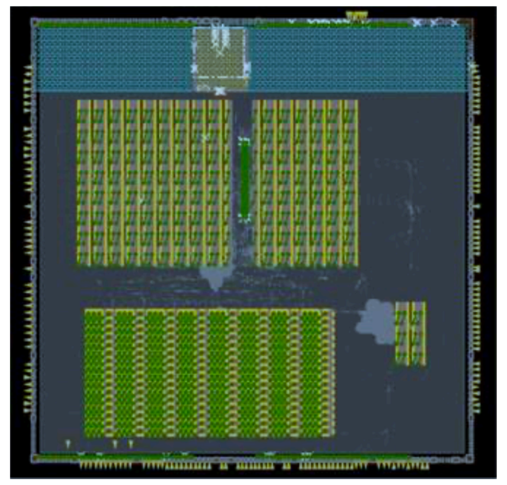
Data Representation Language

Scheduling Language

Compiler

CPUs

GPUs



DSAs

Most of HW retargeting is about changing schedules and data representations